



Network Address Translators (NATs) and NAT Traversal

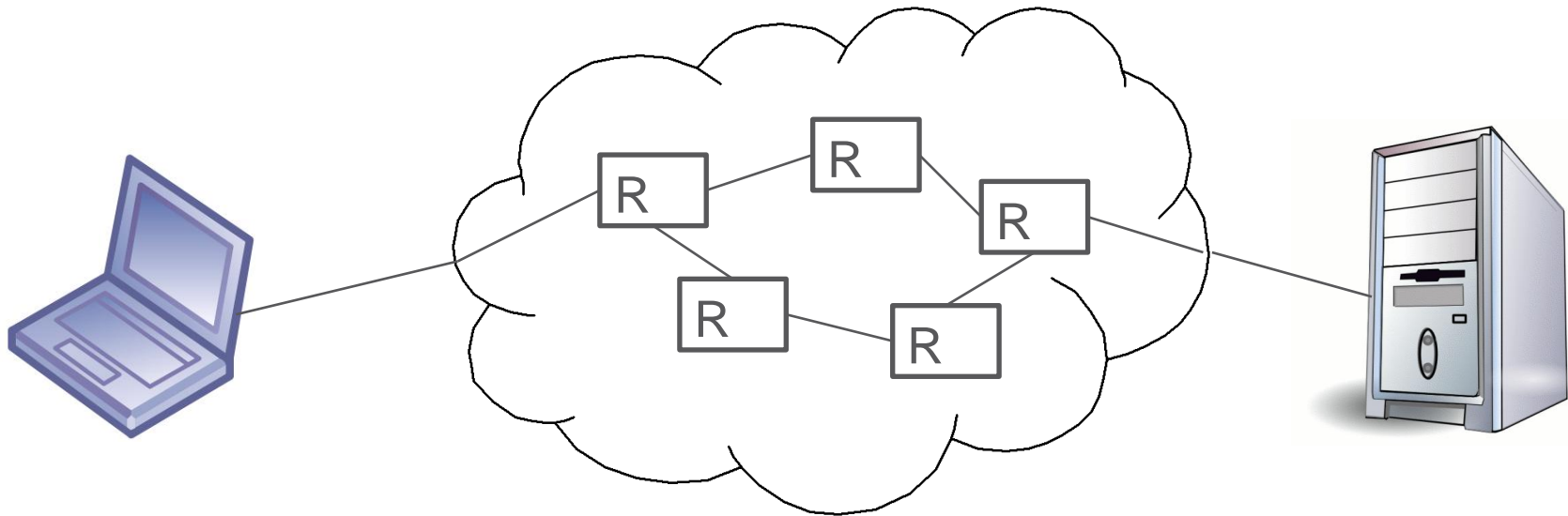
Ari Keränen

Ericsson Research Finland, NomadicLab

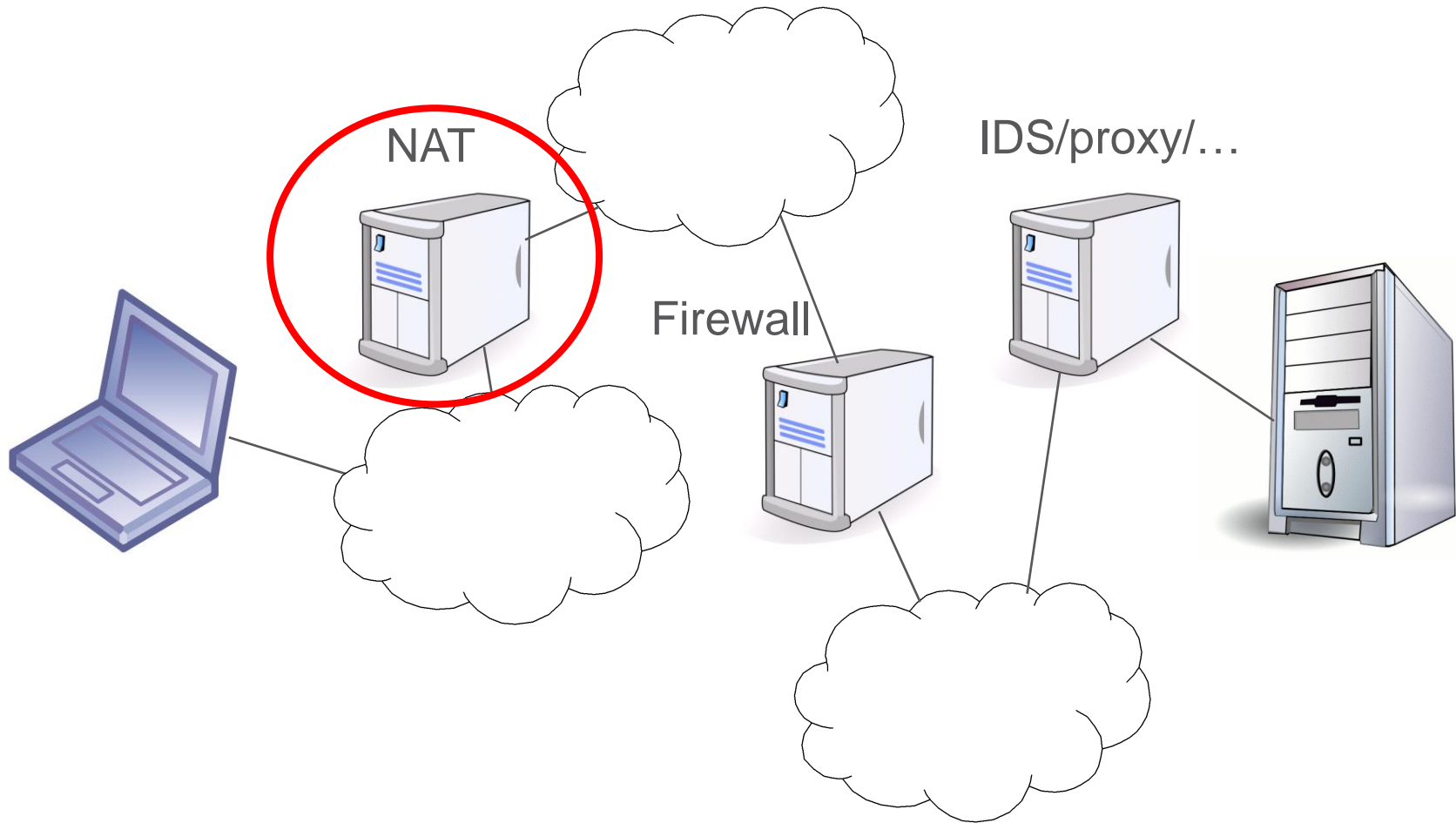
Outline

- › Introduction to NATs
- › NAT Behavior
 - UDP
 - TCP
- › NAT Traversal
 - STUN
 - TURN
 - ICE
 - Others
- › NAT64

Internet Back in the Good Old Days

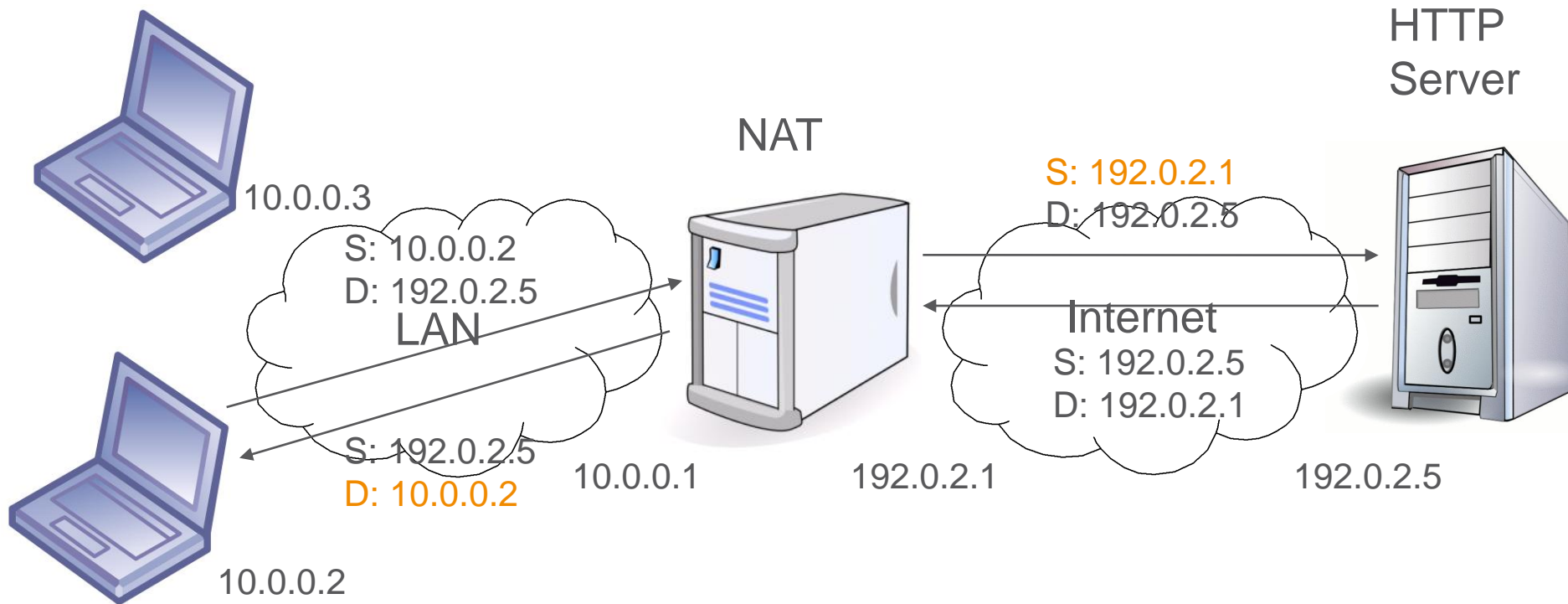


Internet Today (in practice)



Origin of NATs

- › Created to resolve the IPv4 address exhaustion problem
 - Use private address space in the LAN, translate to/from Internet
- › Designed with the web in mind
 - Client/server paradigm



Different Kind of NATs

- › "Basic" Network Address Translator
 - Translates just the IP address in the packets
 - Requires multiple addresses from the NAT
 - › One for each host concurrently communicating with the outside networks
- › Network Address and Port Translator (NAPT)
 - Uses also transport layer (TCP/UDP) ports for multiplexing connections
 - Most of the current NATs are of this type
 - The term "NAT" usually means NAPT
- › NAT64
 - More about this later
- › ...

Side-effects of NATs

- › Hosts behind NATs are not reachable from the public Internet
 - Sometimes used to implement security (should use firewall instead)
 - Breaks peer-to-peer (as opposed to client/server) applications
- › NATs attempt to be transparent
 - Troubleshooting becomes more difficult
- › NATs are a single point of failure
- › NATs may try to change addresses also in the payload (and possibly break application layer protocols)
- › NATs' behavior is not deterministic
 - Difficult to build applications that work through NATs

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IETF NAT Behavior Recommendations

- › RFCs describing how NATs **should** behave
 - RFC 4787: Network Address Translation (NAT) Behavioral Requirements for Unicast UDP
 - RFC 5382: NAT Behavioral Requirements for TCP
- › Classification of current NAT behavior
 - Existing terminology was confusing
 - › Full cone, restricted cone, port restricted cone, and symmetric
- › Recommendations for NAT vendors
 - BEHAVE-compliant NATs are deterministic
- › Lots of NATs implemented before the recommendations
 - Various kind of behavior found in the wild
 - Not all new NATs comply even today

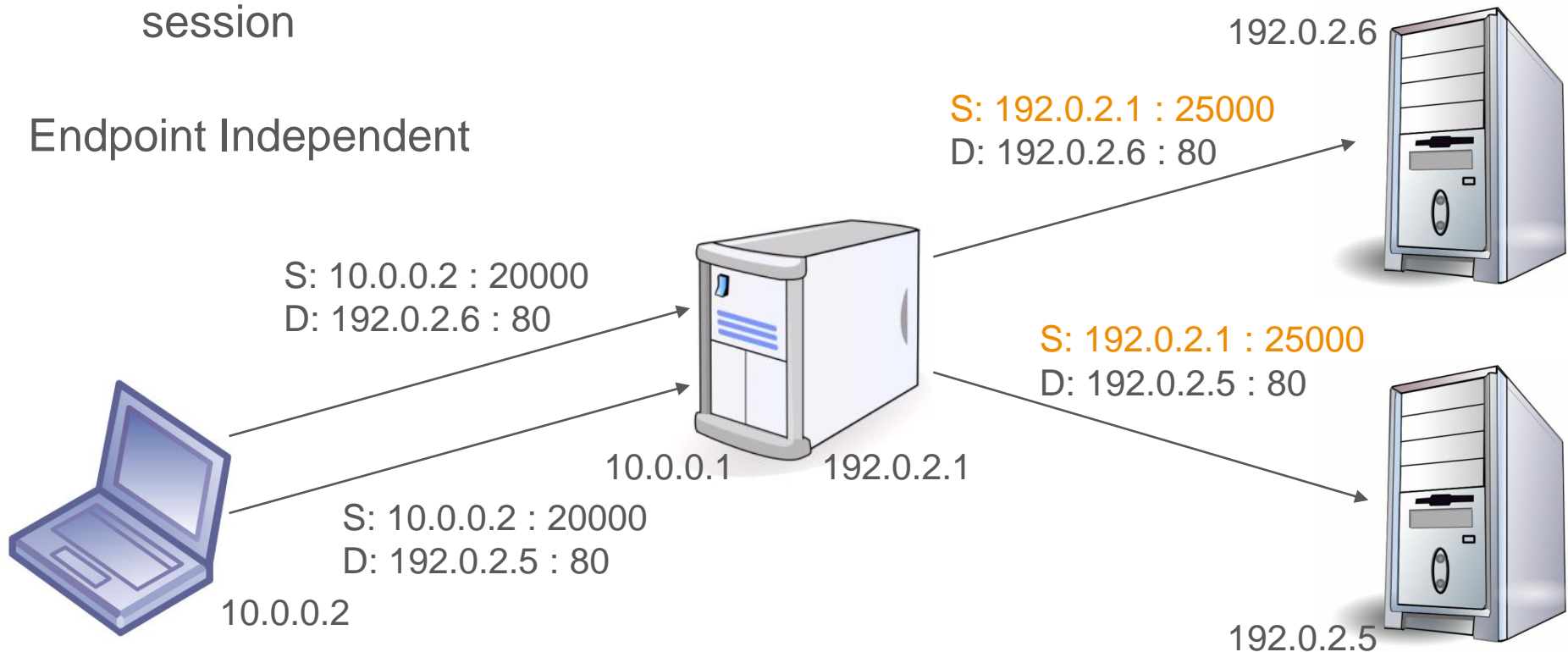
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Mapping Behavior

- › For session originated on the same address and port
 - Endpoint independent: same mapping to different sessions
 - › MUST use it
 - Address dependent: same mapping to sessions to the same host
 - Address and port dependent: a mapping only applies to one session

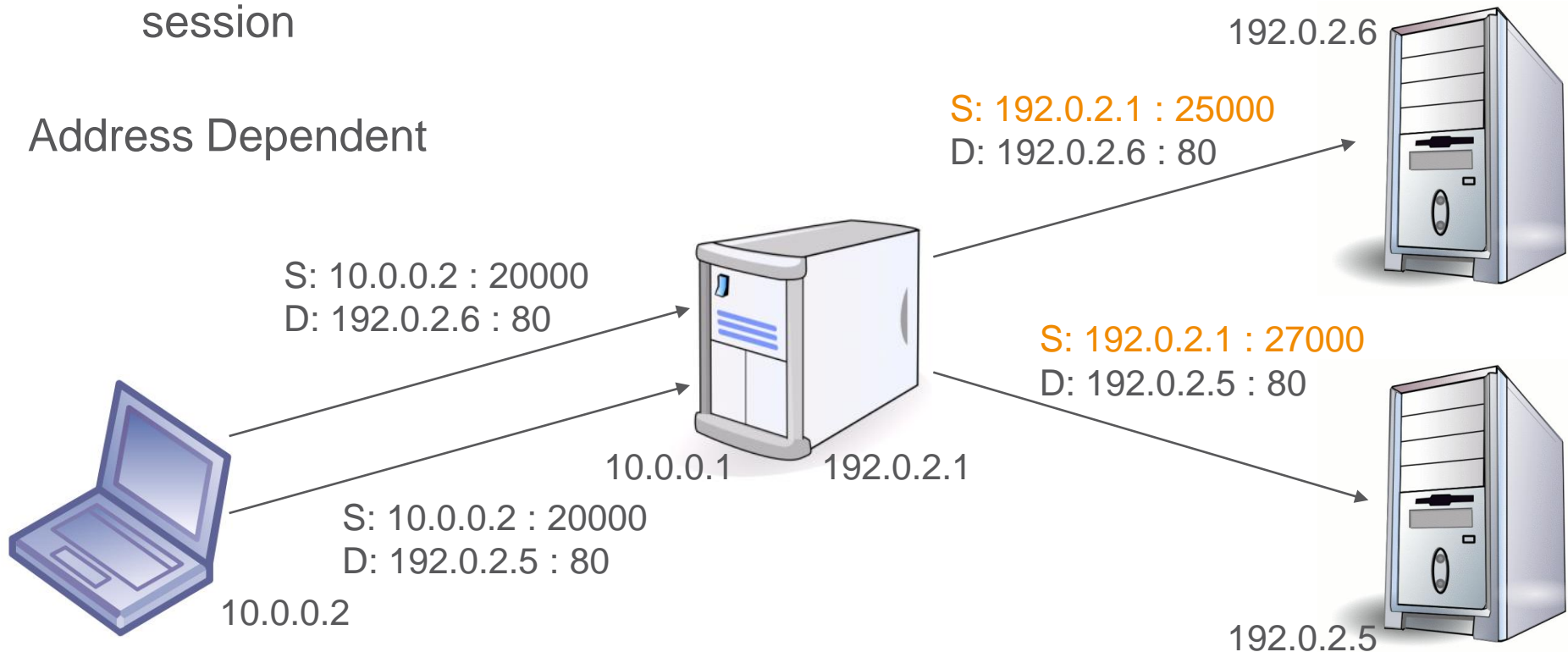
Endpoint Independent



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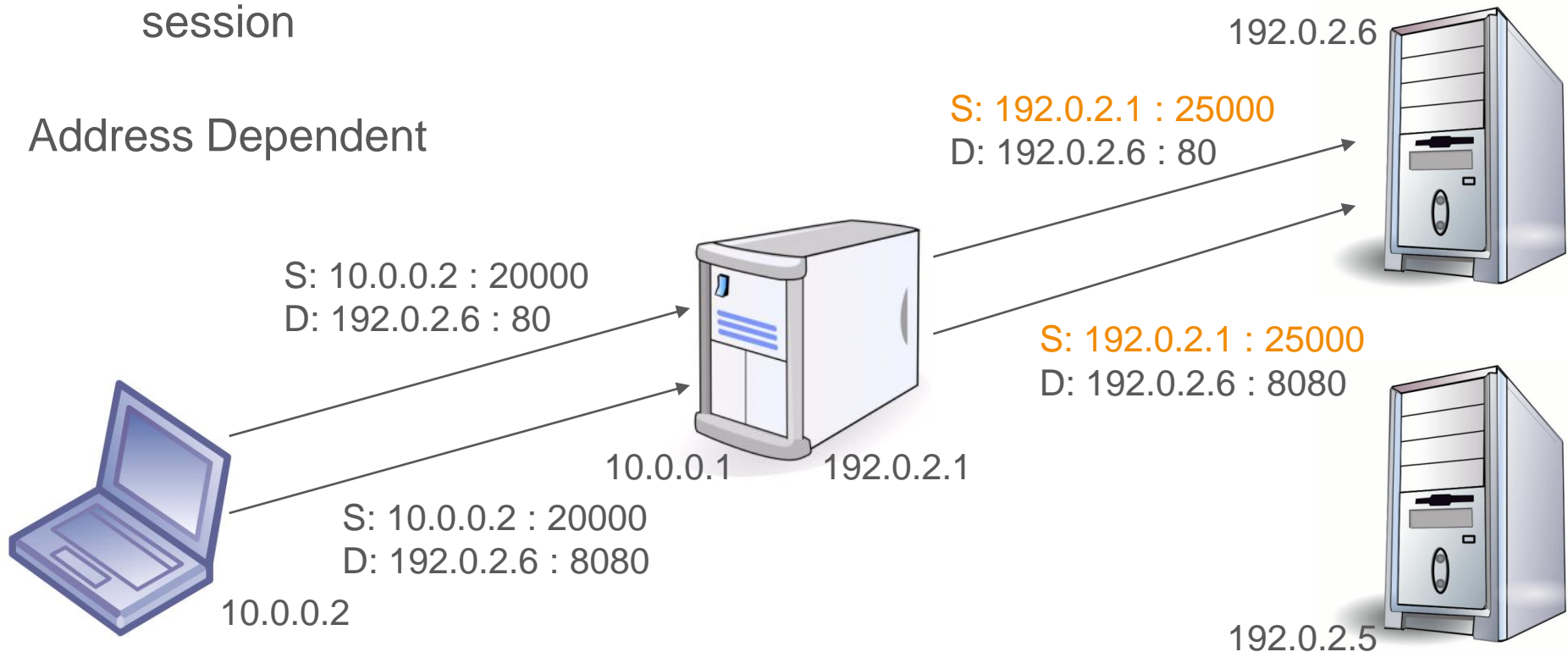
Address Dependent



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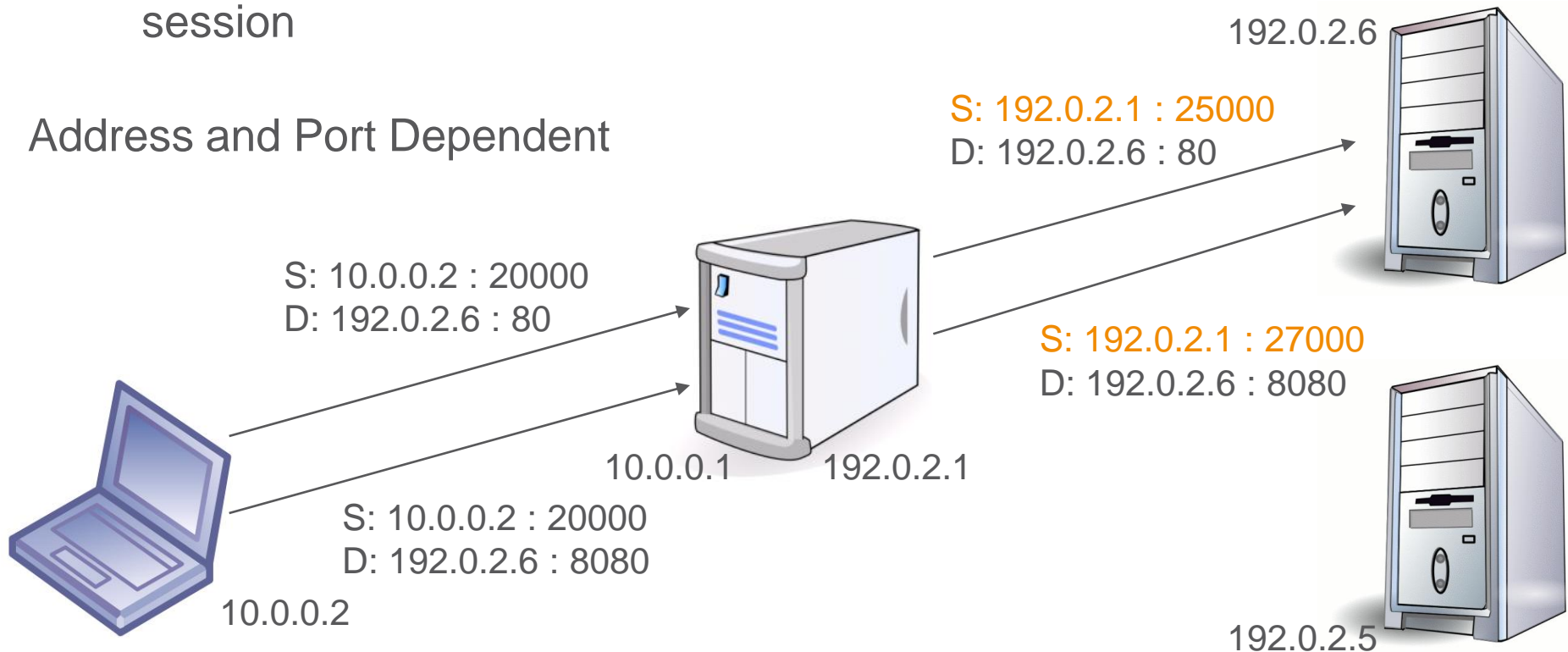
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Address and Port Dependent



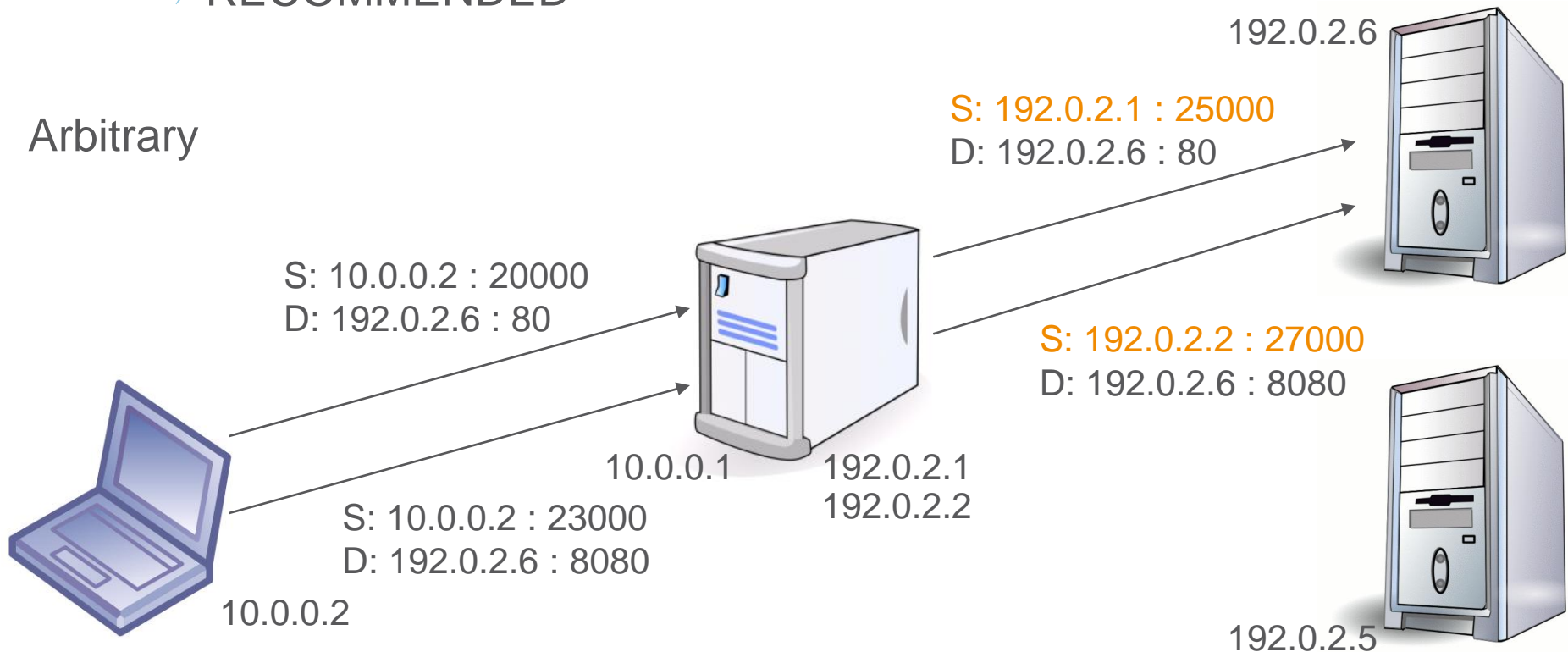
IP Address Pooling Behavior

› NATs with a pool of external IP addresses

- Arbitrary: an endpoint may have simultaneous mappings corresponding to different external IP addresses of the NAT
- Paired: same external IP address of the NAT

› RECOMMENDED

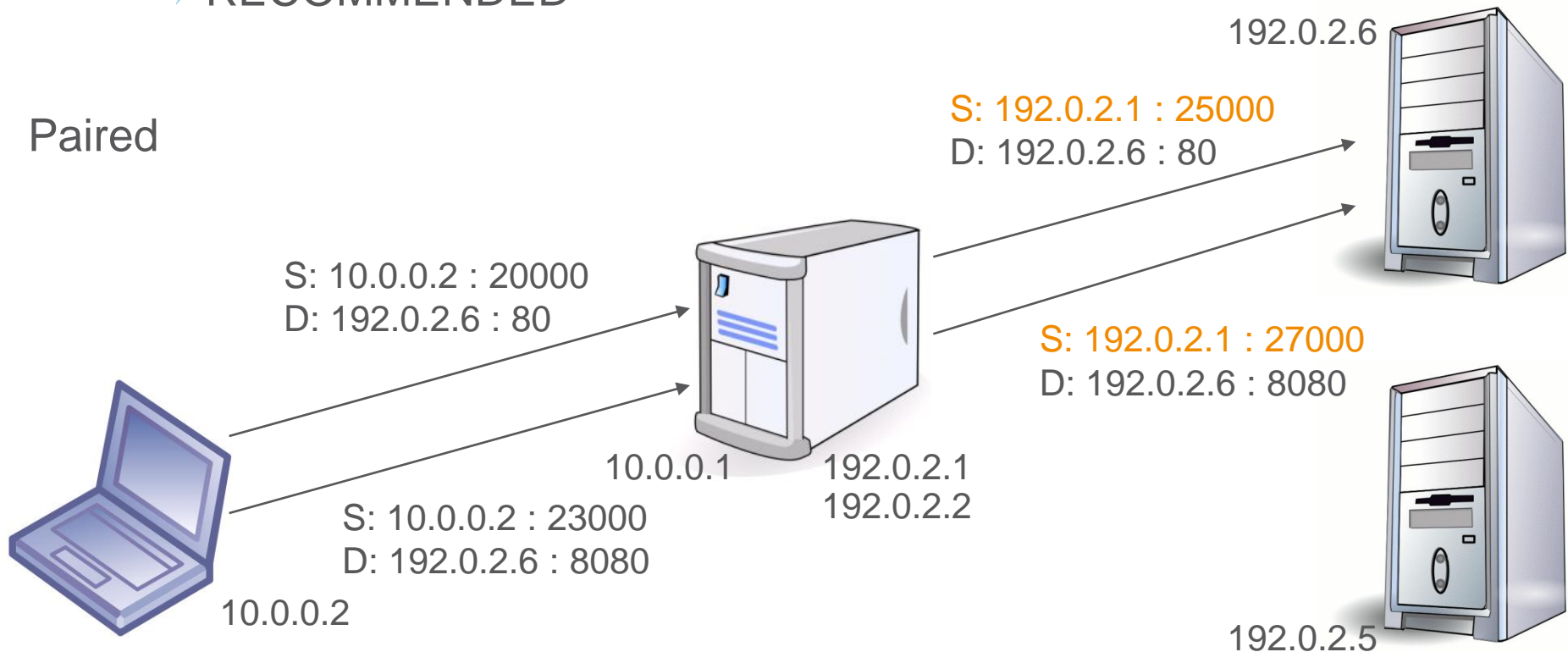
Arbitrary



IP Address Pooling Behavior

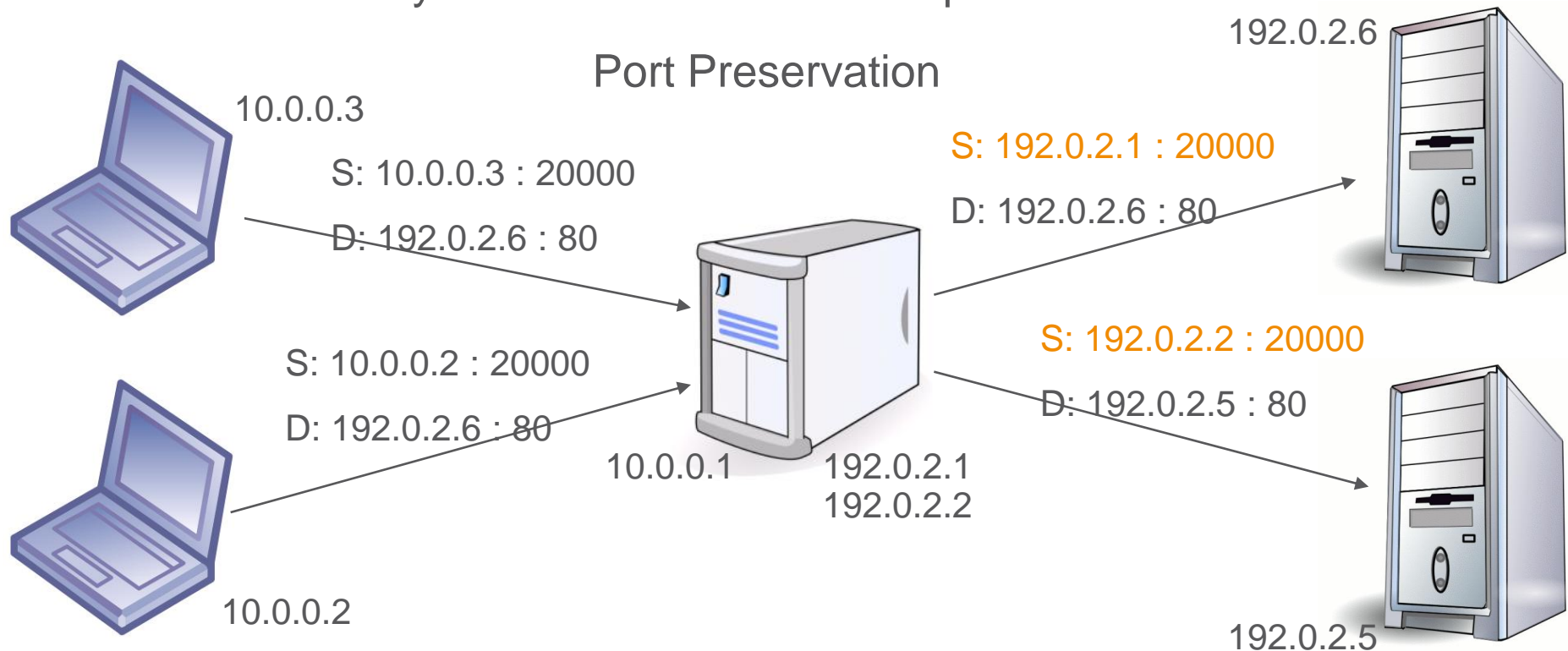
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Paired



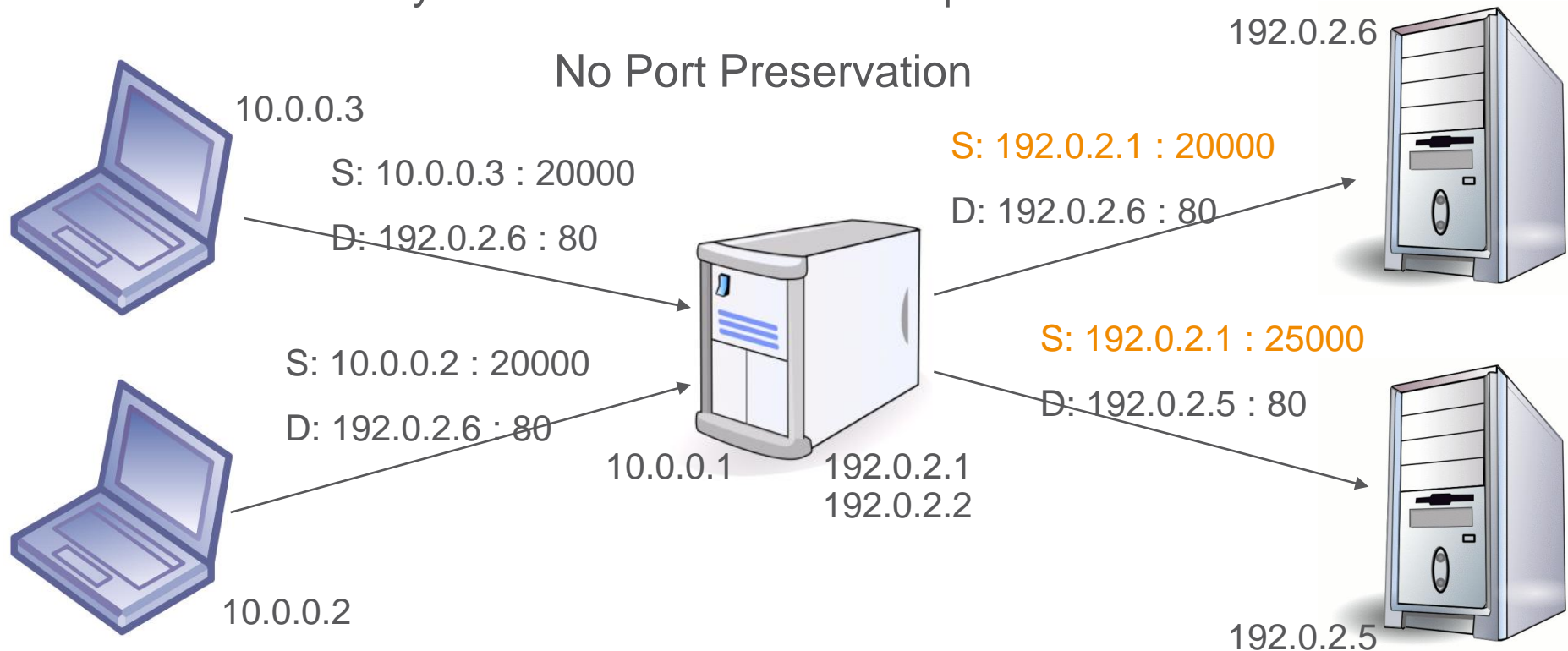
Port Assignment

- › Port preservation: preserves the port as long as there are available IP addresses in the NAT's pool
- › Port overloading: the port is preserved always, even without available IP addresses in the NAT's pool
 - The NAT relays on the source of the response



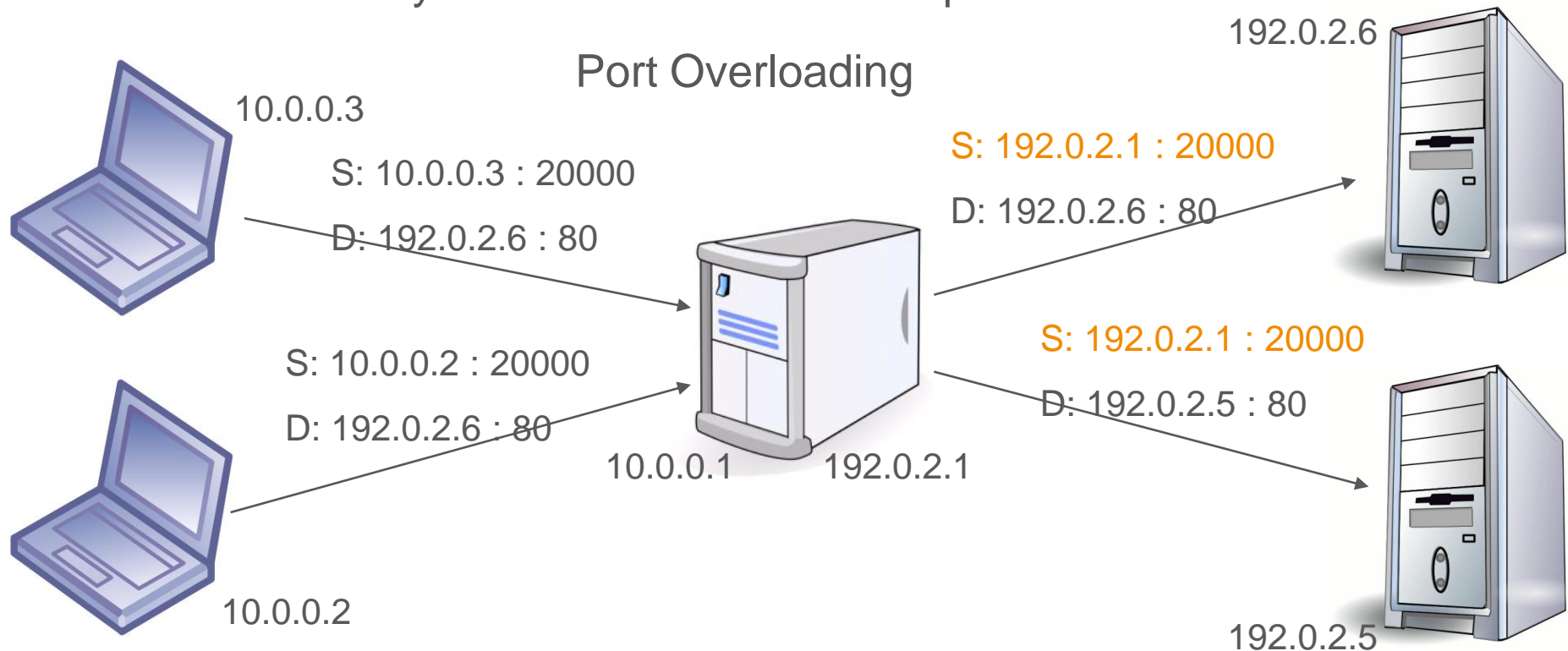
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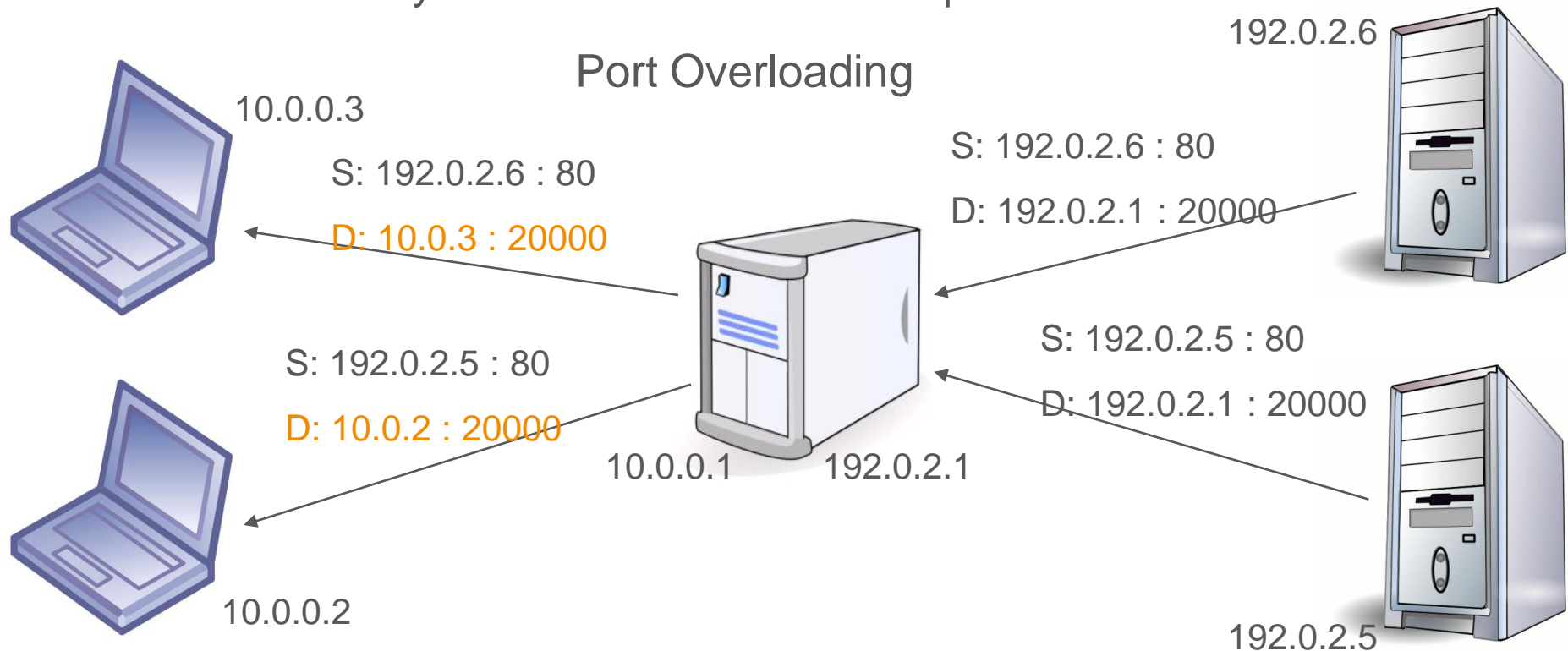
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Port Ranges

- › 1- 1023 Well known
- › 1024 – 49151 Registered
- › 49152 – 65535 Dynamic / Private

- › RECOMMENDED to preserve the following ranges
 - 1 – 1023
 - 1024 – 65535

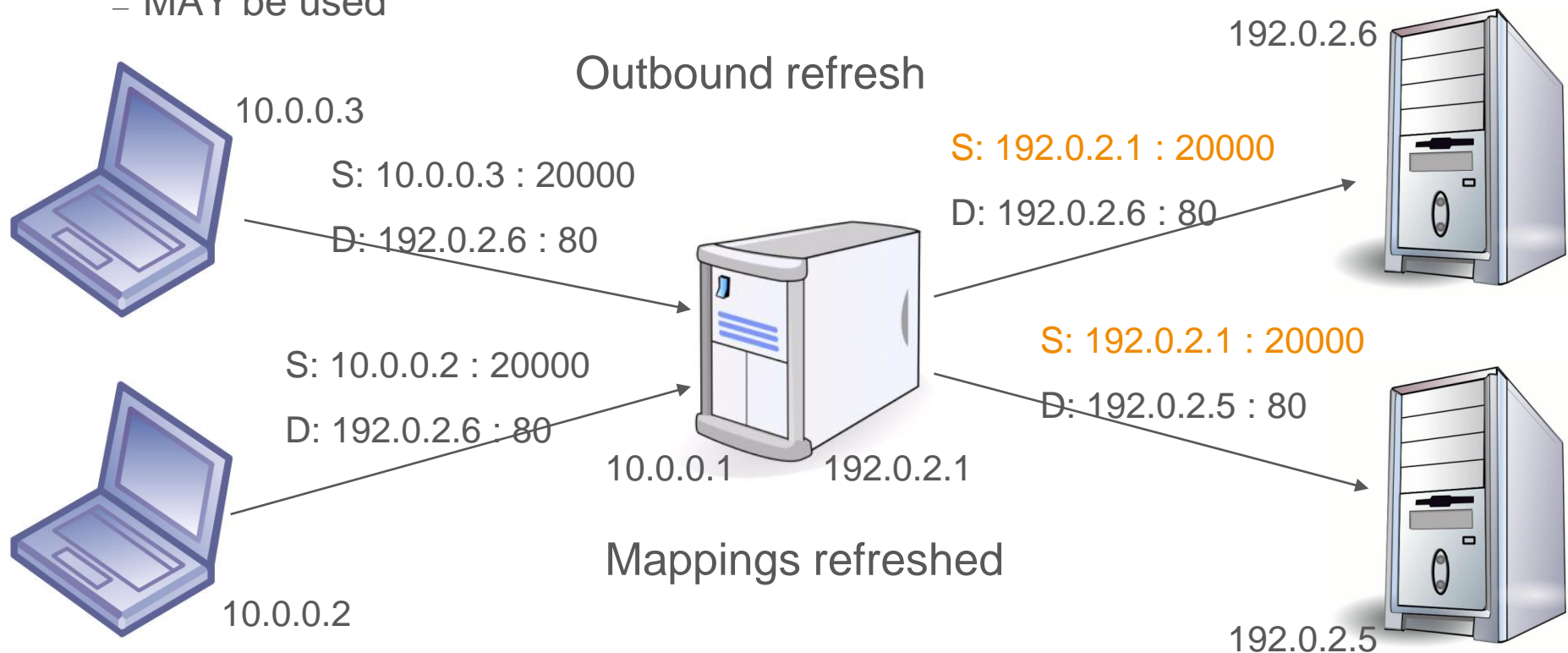
- › Port overloading **MUST NOT** be used
 - Problems when two internal hosts connect to the same external host
- › It is RECOMMENDED that NATs preserve port parity (even/odd)
- › No requirement for port contiguity

Mapping Timeout

- › A NAT UDP mapping **MUST NOT** expire in less than 2 minutes
- › NATs can have application-specific timers
 - Well-known ports
- › It is **RECOMMENDED** to use more than 5 minutes

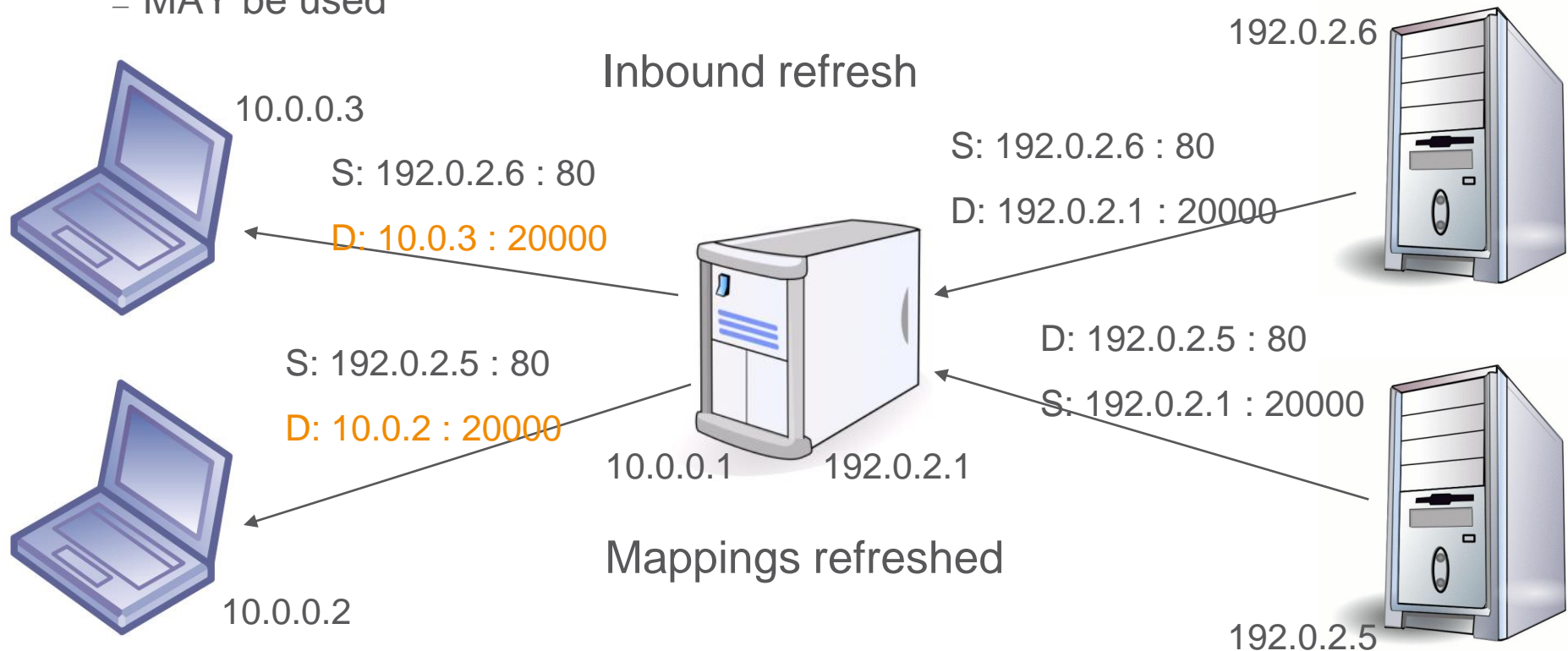
Mapping Refresh

- › NAT outbound refresh: packets from the internal to the external interface
 - MUST be used
- › NAT inbound refresh: packets from the external to the internal interface (attackers may keep the mapping from expiring)
 - MAY be used



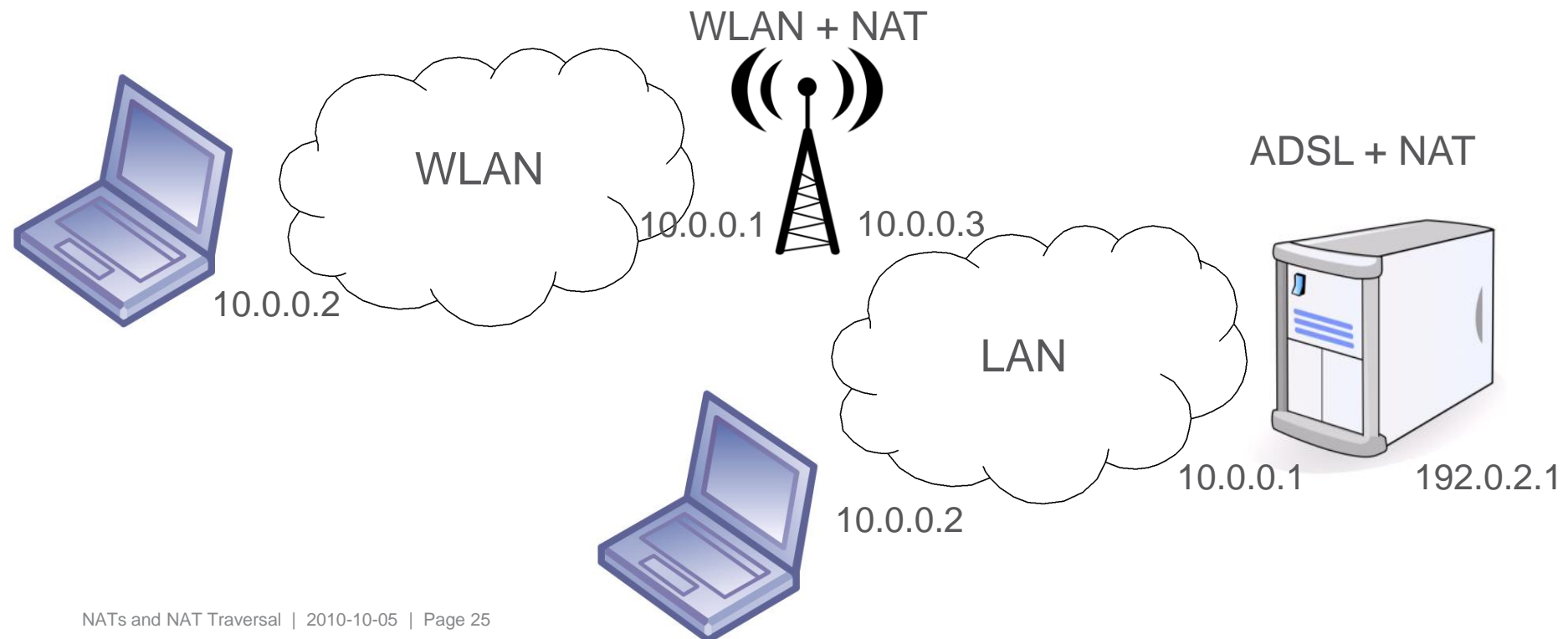
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External Address Spaces

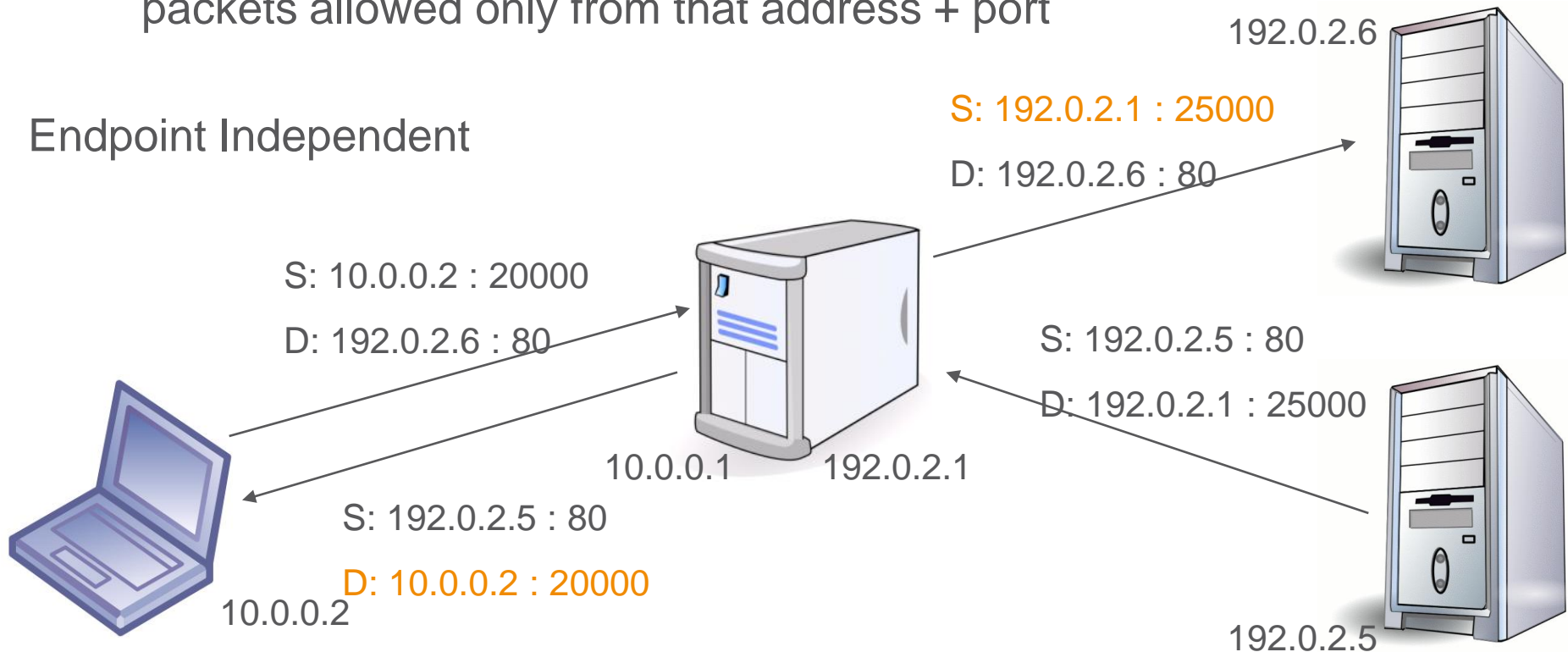
- › NATs MUST be able to handle external address spaces that overlap with the internal address space
 - Internal nodes cannot communicate directly with external nodes that have the same address as another internal node
 - However, they can use STUN techniques



Filtering Behavior

- › Endpoint independent: any packets allowed back
- › Address dependent: external hosts can return packets
- › Address and port dependent
 - Packets sent to an address + port → incoming packets allowed only from that address + port

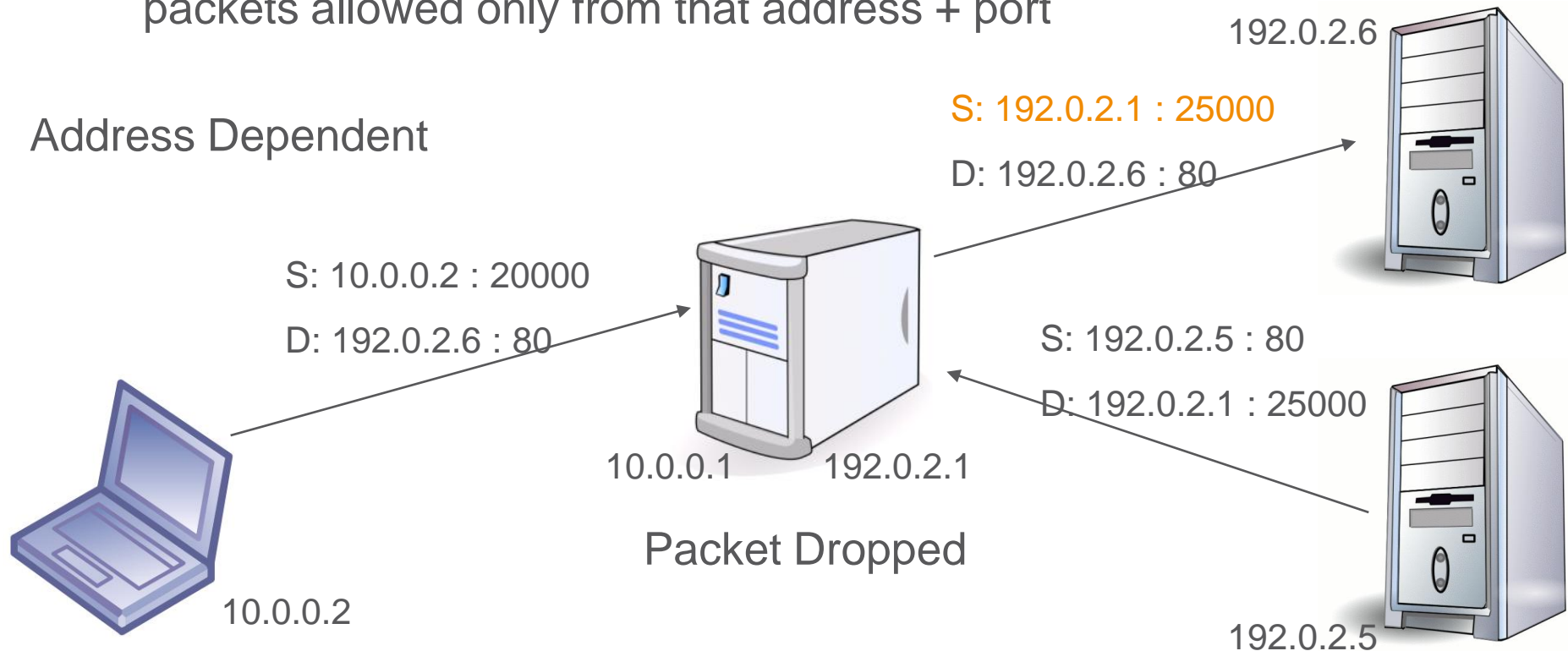
Endpoint Independent



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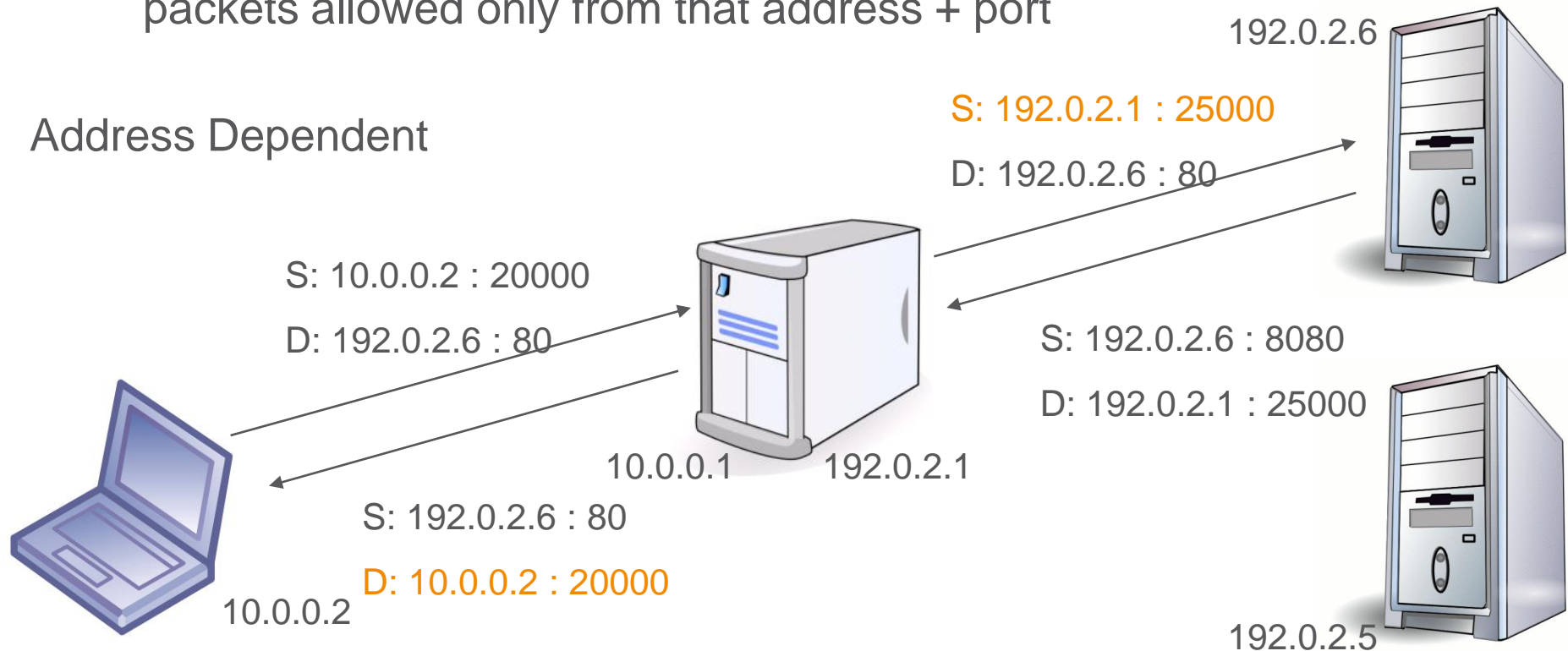
Address Dependent



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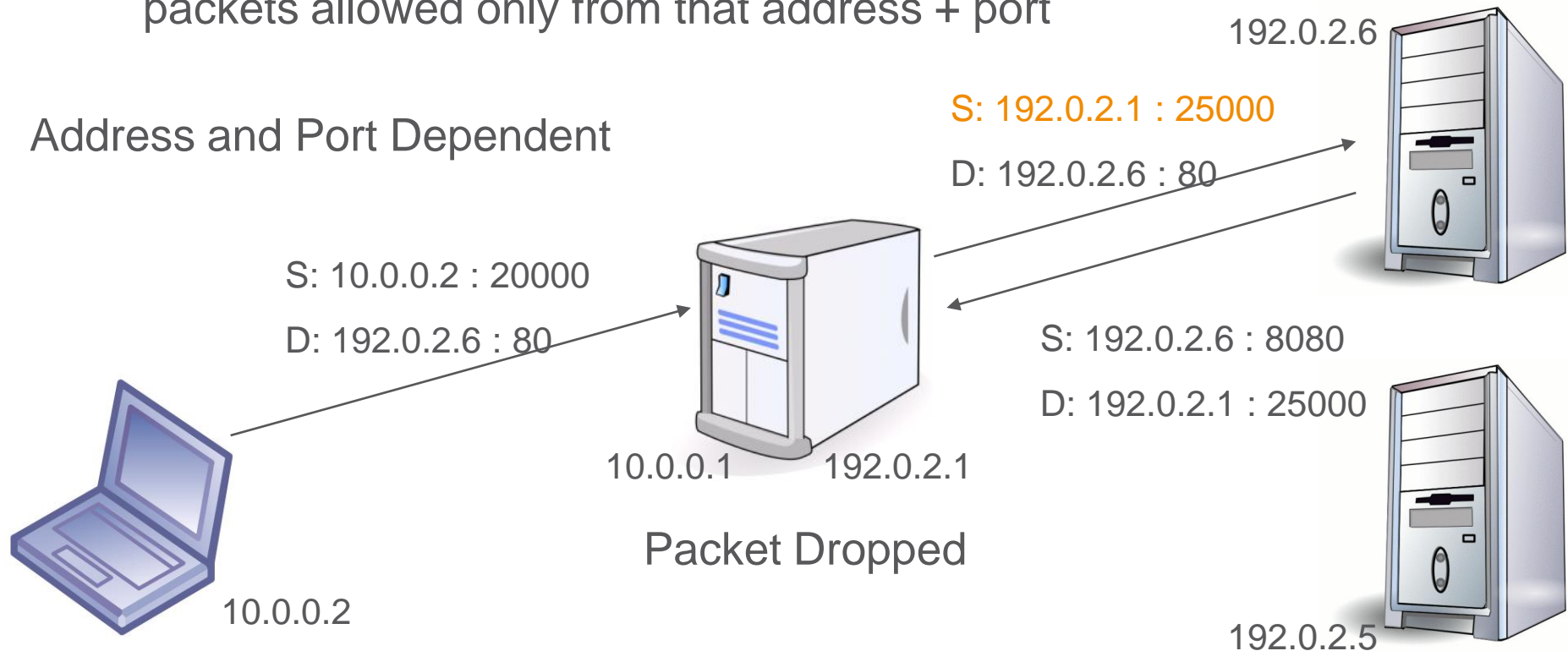
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Address and Port Dependent

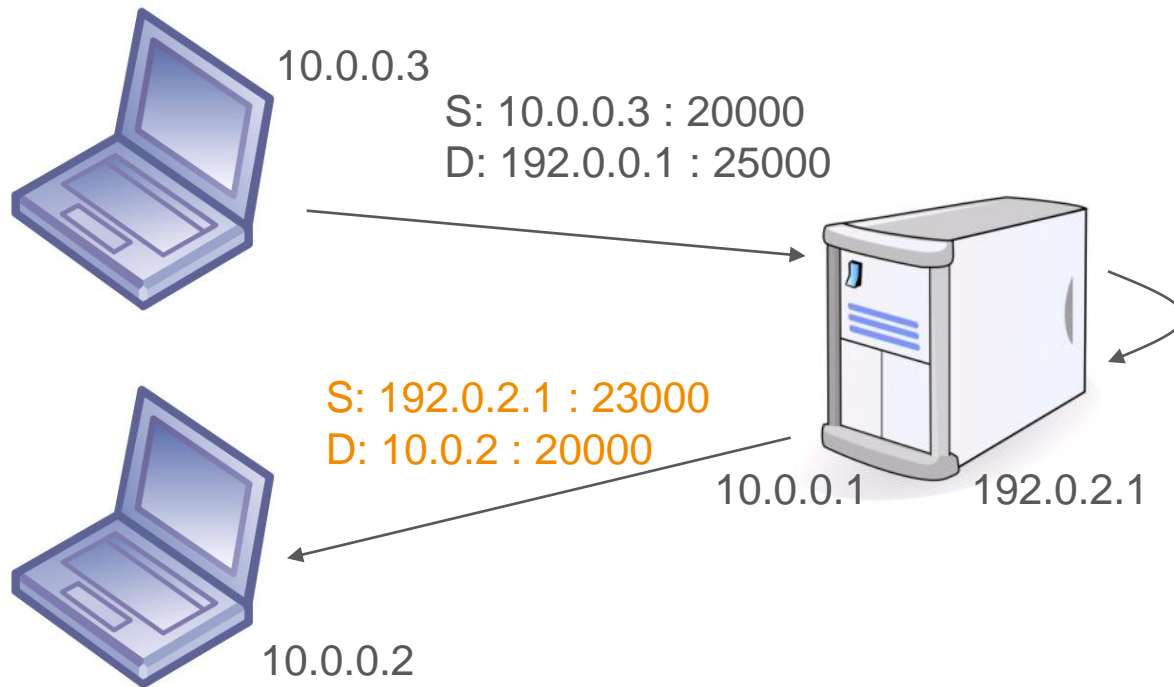


Filtering Behavior

- › Endpoint independent filtering is RECOMMENDED
 - Opens up ports for attackers
- › If a more stringent filtering is required
 - Address dependent filtering is RECOMMENDED

Hairpinning

- › Internal hosts communicate using external addresses
 - MUST be supported



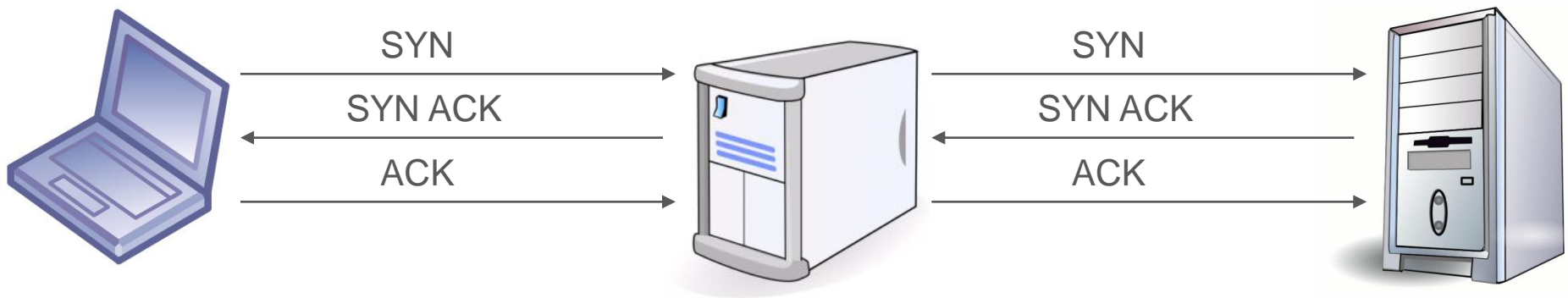
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TCP Connection Establishment

- › Three-way handshake
 - MUST be supported
- › Simultaneous open
 - MUST be supported

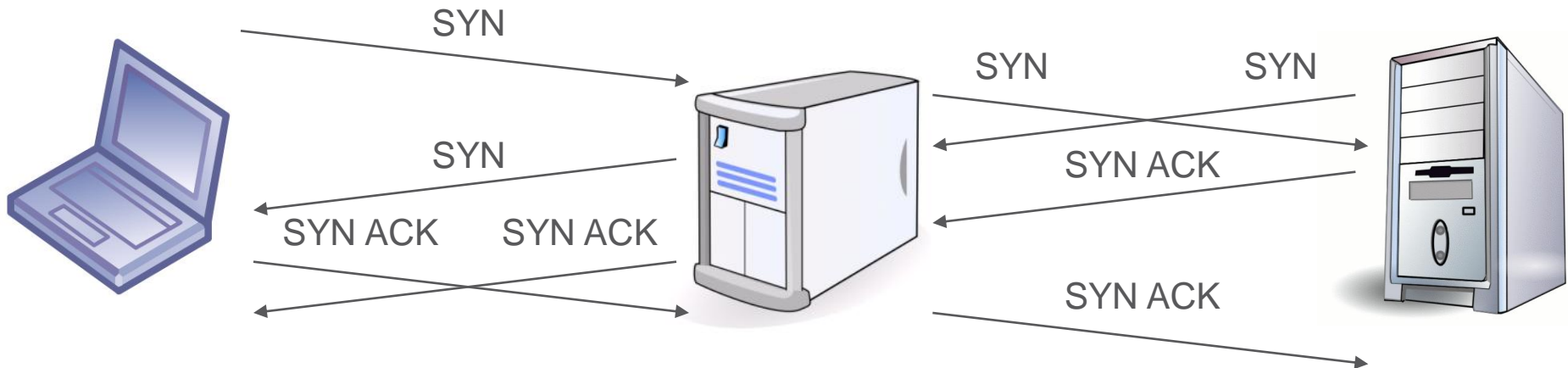
Three-way Handshake



TCP Connection Establishment

- › Three-way handshake
 - MUST be supported
- › Simultaneous open
 - MUST be supported

Simultaneous Open



NAT Session Timeout

- › Established connections
 - MUST NOT be less than 2 hours and 4 minutes
 - By default TCP keepalives are sent every 2 hours
- › Partially opened or partially closed connections
 - MUST NOT be less than 4 minutes
- › TIME_WAIT timeout not specified

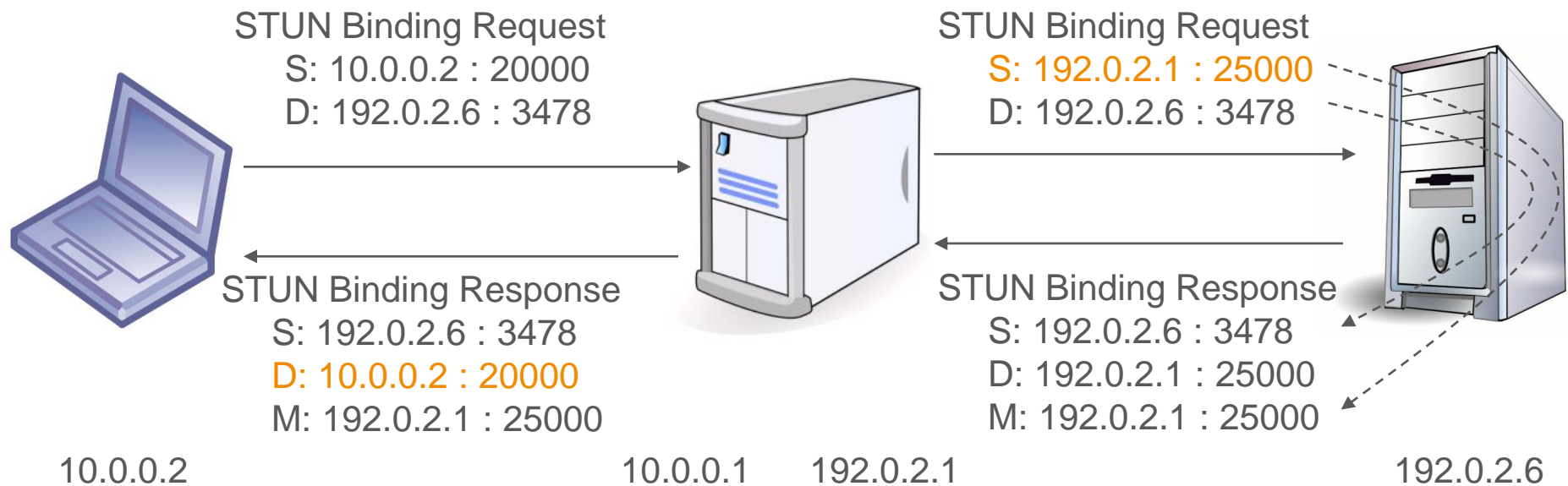
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STUN

- › Session Traversal Utilities for NAT (RFC 5389)
- › Originally a protocol between endpoints and “reflectors”
- › Revised specification defines usages
 - Binding discovery using STUN servers
 - NAT keepalives
 - Authentication (short-term password and long term credentials)
- › TLV encoded
- › Can run on UDP, TCP, or TLS/TCP
- › STUN server discovered using DNS SRV
- › Transactions
 - Request/response
 - Indications (not delivered reliably)
- › Can be multiplexed with other protocols
 - Two first bits are zeros
 - Magic cookie
 - FINGERPRINT attribute

Binding Discovery



XOR-MAPPED-ADDRESS

- › Some NATs inspect packets and translate IP addresses known to them
 - Try to be smart and “fix” the application layer protocol
- › The mapped address is obfuscated in the response so that NAT does not recognize it
 - Simple XOR operation

Outline

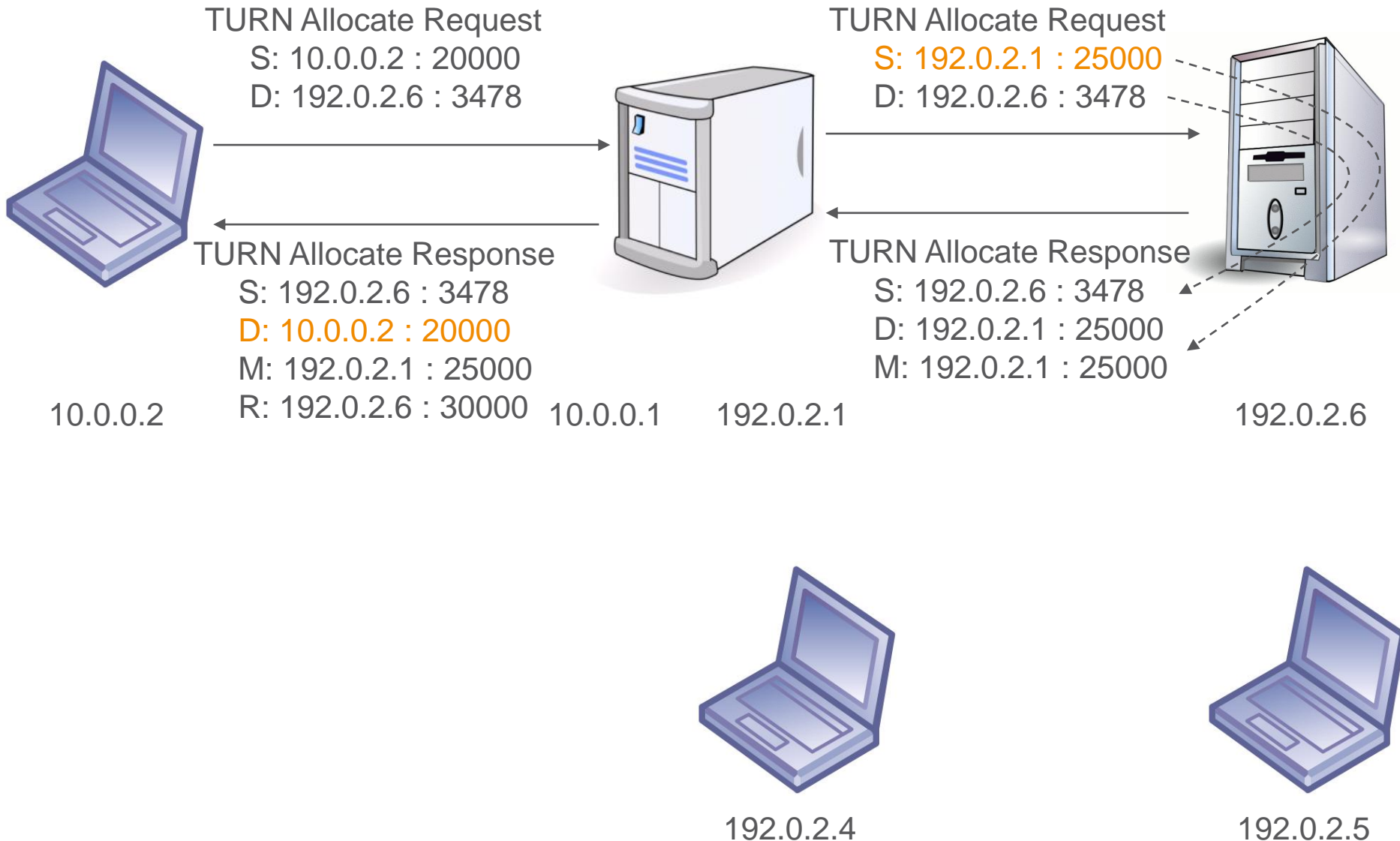
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TURN

- › Traversal Using Relays around NAT: Relay Extensions to Session Traversal Utilities for NAT (RFC 5766)
- › Allocate request / response
 - Allocate an external “relayed” address at the relay
 - Responses carry the mapped and the relayed address
- › Send and Data indication
 - STUN messages containing relayed data
 - Send data to a remote endpoint through the relay
 - Data received from remote endpoints through the relay
- › Channels
 - Send and receive relayed data with minimalistic (32-bit) header
- › Permissions

Relay Operations

R: 192.0.2.6 : 30000

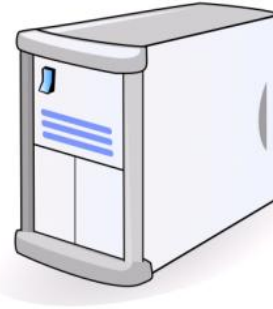


Relay Operations

R: 192.0.2.6 : 30000



10.0.0.2



10.0.0.1

192.0.2.1



192.0.2.6

Packet Dropped

The client needs to set a permission in the relay in order to receive data through it

Equivalent to a NAT with:

Address dependent filtering policy

Endpoint independent mapping

S: 192.0.2.4 : 27000

D: 192.0.2.6 : 30000



192.0.2.4



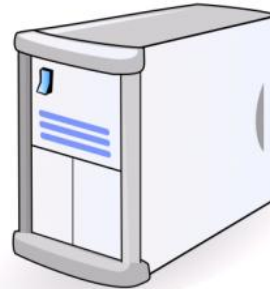
192.0.2.5

Relay Operations

R: 192.0.2.6 : 30000



10.0.0.2



10.0.0.1

192.0.2.1



192.0.2.6

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D: 192.0.2.6 : 30000

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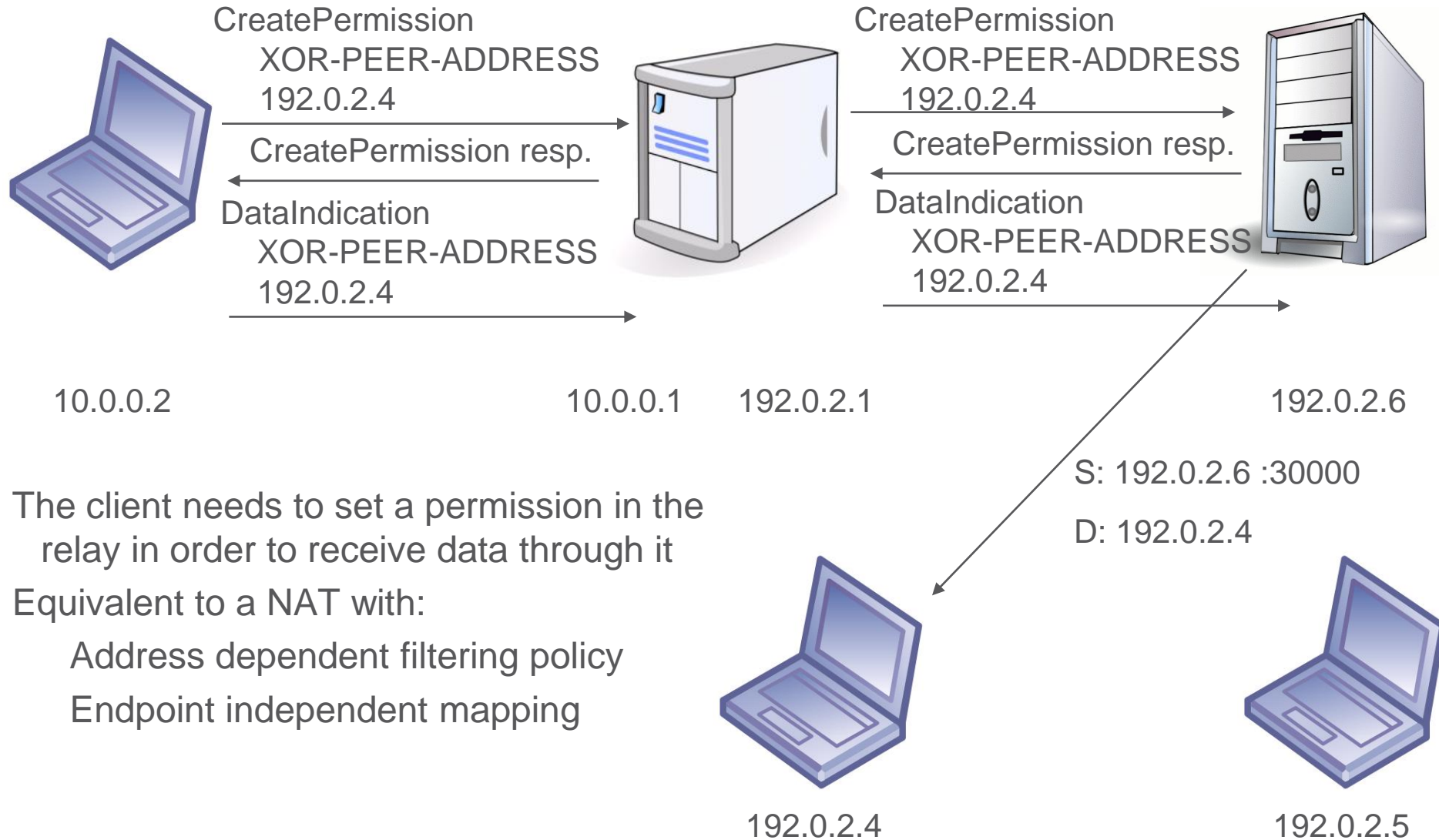
192.0.2.4



192.0.2.5

Relay Operations

R: 192.0.2.6 : 30000



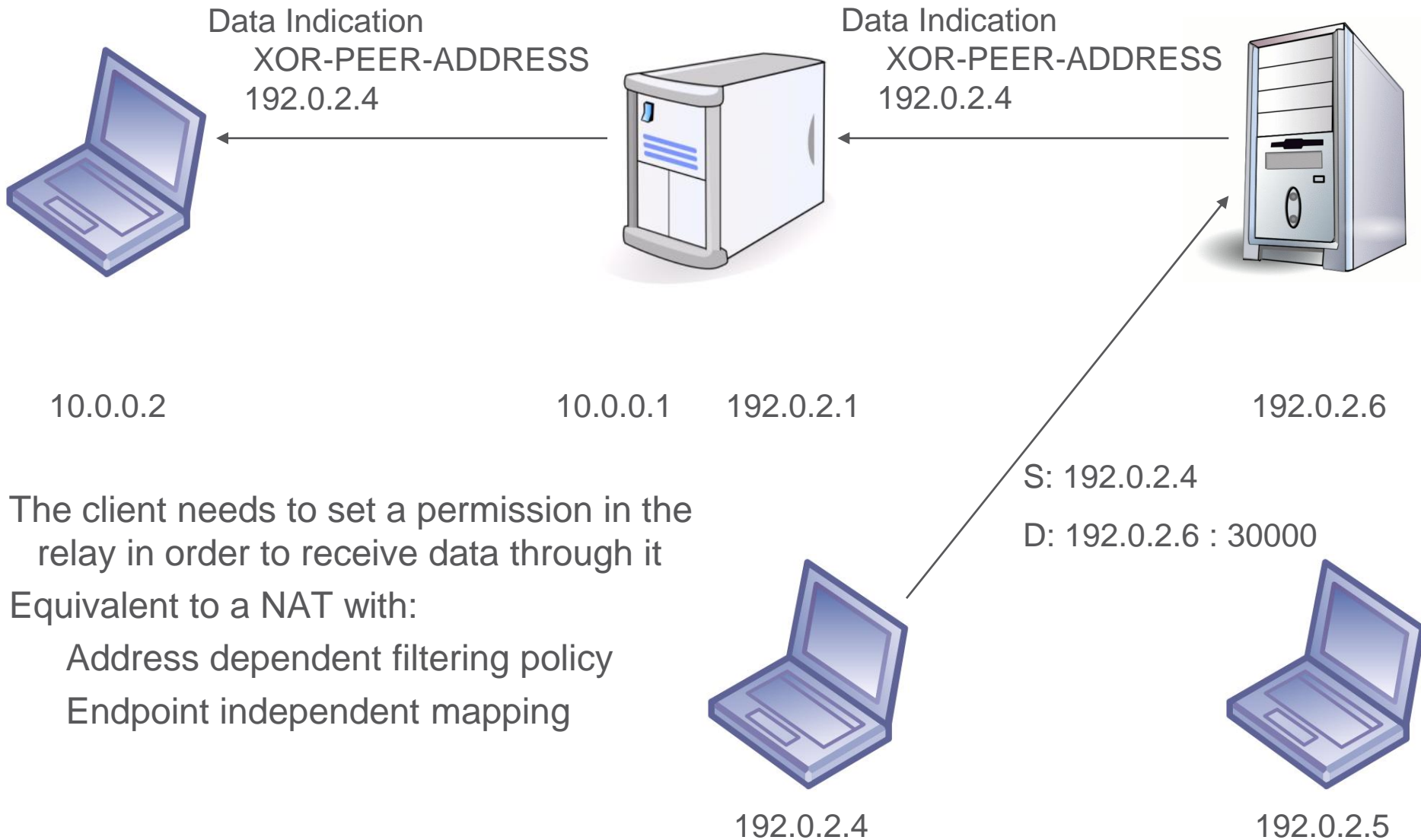
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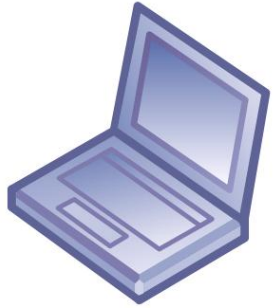
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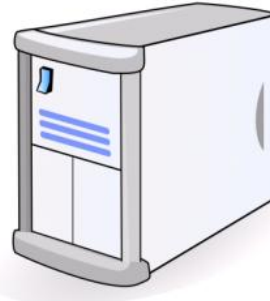


Relay Operations

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10.0.0.2



10.0.0.1

192.0.2.1



192.0.2.6

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192.0.2.4



192.0.2.5

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ICE

- › Interactive Connectivity Establishment : A Protocol for Network Address Translator Traversal for Offer/Answer Protocols (RFC 5245)
- › Uses and extends STUN and TURN protocols
- › Overall procedure:
 - Endpoints gather all the addresses they can
 - › Using e.g. STUN and/or TURN
 - Addresses (candidates) are exchanged with the peer
 - Connectivity checks are run between the candidates
 - The highest priority candidate pair that works is selected for use

Gathering Addresses

- › Address types
 - Host candidates
 - Server-reflexive candidates
 - Relayed candidates
 - Peer-reflexive candidates
- › Duplicated addresses are removed
- › Foundation: used to freeze addresses (related to connectivity checks)
 - Same type
 - Bases with the same IP address
 - Same STUN server

Prioritizing Addresses

$$\text{Priority} = 2^{24} (\text{type preference}) + 2^8 (\text{local preference}) + 2 (256 - \text{component ID})$$

- › Type preference [0-126]: preference for the type of candidate (e.g., server reflexive)
- › Local preference [0-65535]: preference for the interface the candidate was obtained from (e.g., multihomed hosts)
- › Component ID [1-256]: for media with multiple components (e.g., RTP and RTCP)

Connectivity Checks

- › Five states for a pair:
 - Waiting, in progress, succeeded, failed, frozen
- › Periodic checks and triggered checks
 - Periodic checks performed in priority order
 - Incoming check may cause a triggered check
- › Connectivity is checked with STUN Binding Requests
 - Carry a concatenation of user names and the remote password

ICE Roles

› Controlling agent

- Agent that generates the initial offer
- Selects which pair to eventually use
 - › Implementation specific stopping criteria
 - › USE-CANDIDATE attribute

› Controlled agent

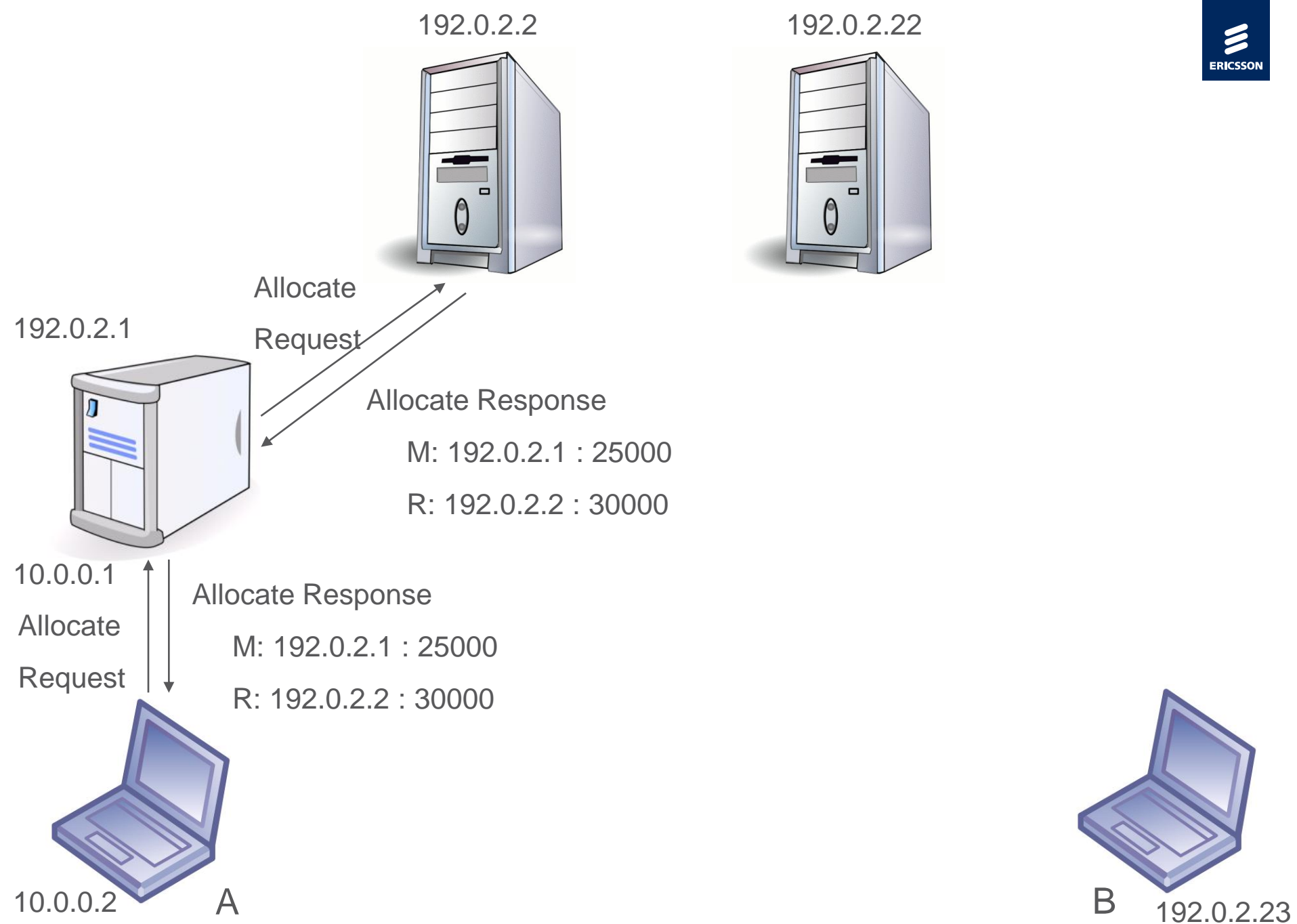
- Generates checks and responds to them like the controlling agent
- Waits for the controlling agent to decide which candidate to use

› ICE lite agents

- Know they are not behind a NAT
 - › e.g., PSTN gateways, conferencing servers
- Always in controlled role
- Just respond to checks

ICE Example (1)

- › One endpoint is behind a NAT
- › One endpoint has a public IP address
- › Endpoints use TURN servers
 - Permission setting is omitted from the examples for brevity



Host candidate:
10.0.0.2 : 20000
Server reflexive:
192.0.2.1 : 25000
Relayed:
192.0.2.2 : 30000

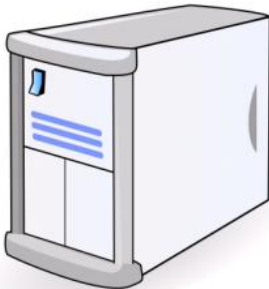
192.0.2.2



192.0.2.22



192.0.2.1



10.0.0.1

Allocate Response

M: 192.0.2.1 : 25000
R: 192.0.2.2 : 30000

INVITE (offer)



10.0.0.2



192.0.2.23

Host candidate:
10.0.0.2 : 20000
Server reflexive:
192.0.2.1 : 25000
Relayed:
192.0.2.2 : 30000

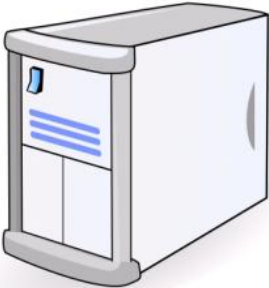
192.0.2.2



192.0.2.22



192.0.2.1



10.0.0.1



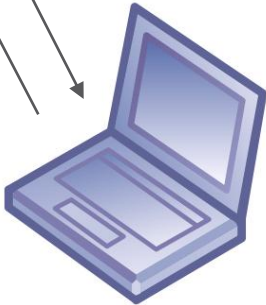
10.0.0.2

Allocate
Request

Allocate Response

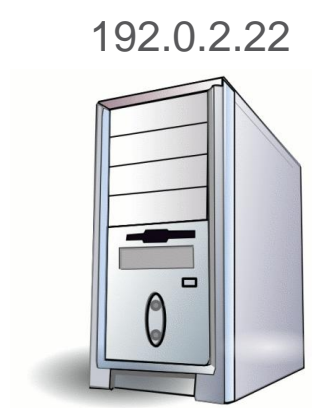
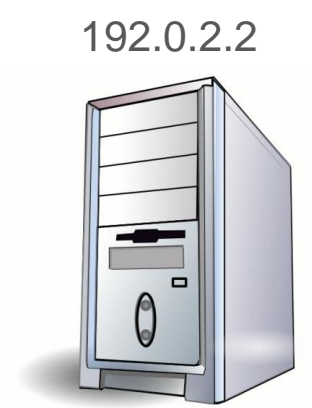
M: 192.0.2.23 : 35000

R: 192.0.2.22 : 45000



192.0.2.23

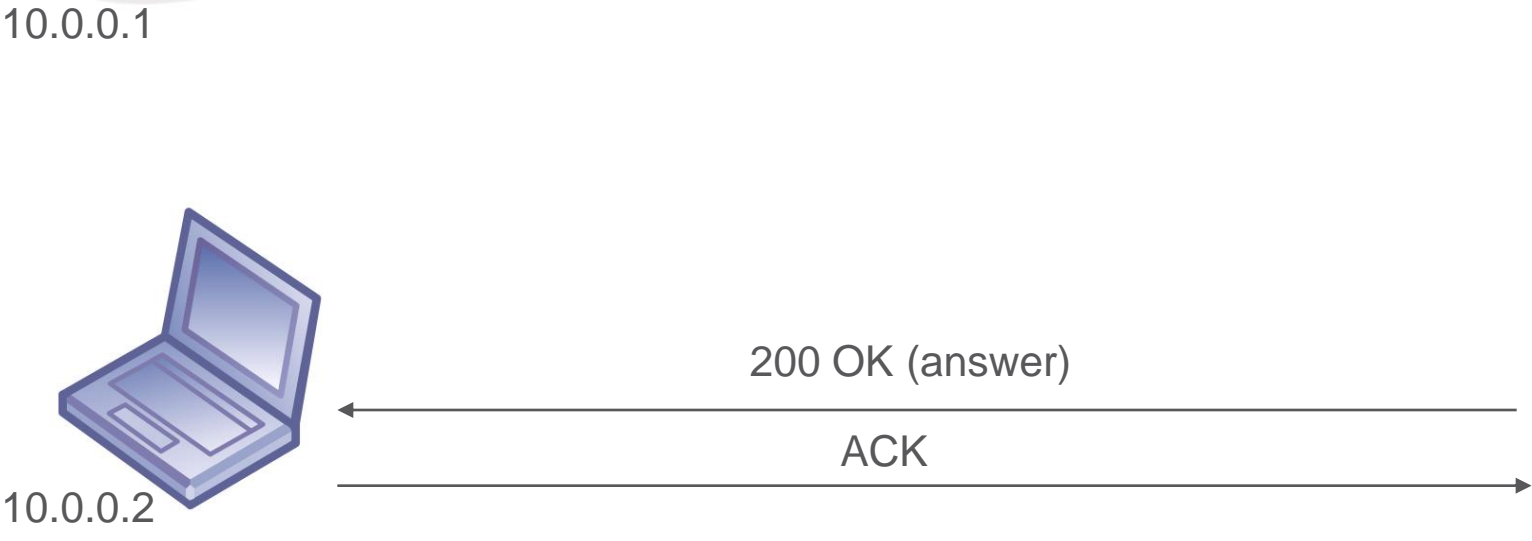
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Relayed:
192.0.2.2 : 30000



Host candidate:
~~192.0.2.23 : 35000~~
~~Server reflexive:~~
~~192.0.2.22 : 45000~~
Relayed:
192.0.2.22 : 45000



Allocate Response
M: 192.0.2.23 : 35000
R: 192.0.2.22 : 45000



192.0.2.2



192.0.2.22



Host candidate:
192.0.2.23 : 35000
Relayed:
192.0.2.22 : 45000

Host candidate:
10.0.0.2 : 20000
Server reflexive:
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Relayed:
192.0.2.2 : 30000

192.0.2.1



10.0.0.1
Binding
Request

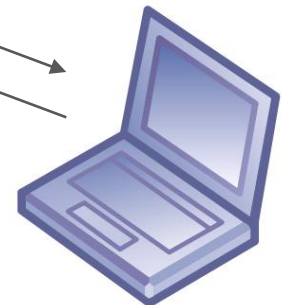
Binding Response
M: 192.0.2.1 : 25000



10.0.0.2

Binding
Request

Binding Response
M: 192.0.2.1 : 25000



192.0.2.23

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 10.0.0.2 : 20000
 Server reflexive:
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 Relayed:
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 Relayed:
 192.0.2.22 : 45000

192.0.2.1



10.0.0.1
 Binding
 Request

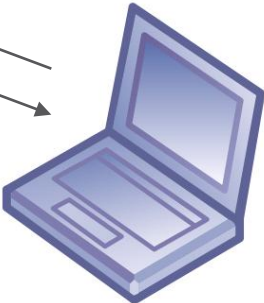
Binding Response
 M: 192.0.2.23 : 35000



10.0.0.2

Binding
 Request

Binding Response
 M: 192.0.2.23 : 35000



192.0.2.23

192.0.2.2



192.0.2.22



Host candidate:
192.0.2.23 : 35000
Relayed:
192.0.2.22 : 45000

Host candidate:
10.0.0.2 : 20000
Server reflexive:
192.0.2.1 : 25000
Relayed:
192.0.2.2 : 30000

192.0.2.1



10.0.0.1
Binding
Request

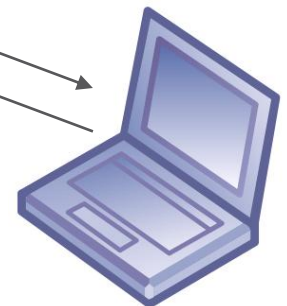


10.0.0.2

Binding Response
M: 192.0.2.23 : 35000

Binding Request
USE-CANDIDATE

Binding Response
M: 192.0.2.23 : 35000



192.0.2.23

192.0.2.2



192.0.2.22



Host candidate:

192.0.2.23 : 35000

Relayed:

192.0.2.22 : 45000

Host candidate:

10.0.0.2 : 20000

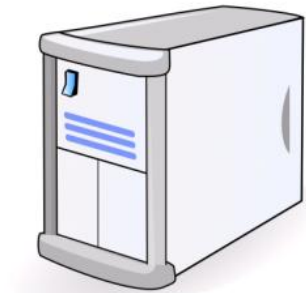
Server reflexive:

192.0.2.1 : 25000

Relayed:

192.0.2.2 : 30000

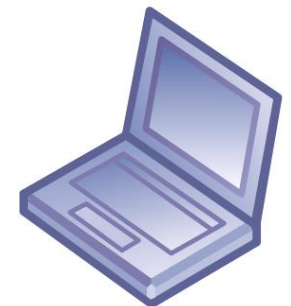
192.0.2.1



10.0.0.1



10.0.0.2



192.0.2.23

INVITE (offer)

200 OK (answer)

ACK

192.0.2.2



192.0.2.22



Host candidate:

192.0.2.23 : 35000

Relayed:

192.0.2.22 : 45000

Host candidate:

10.0.0.2 : 20000

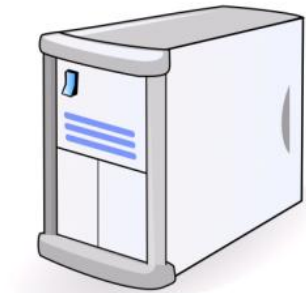
Server reflexive:

192.0.2.1 : 25000

Relayed:

192.0.2.2 : 30000

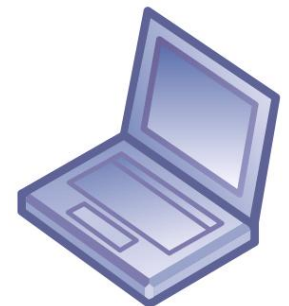
192.0.2.1



10.0.0.1



10.0.0.2

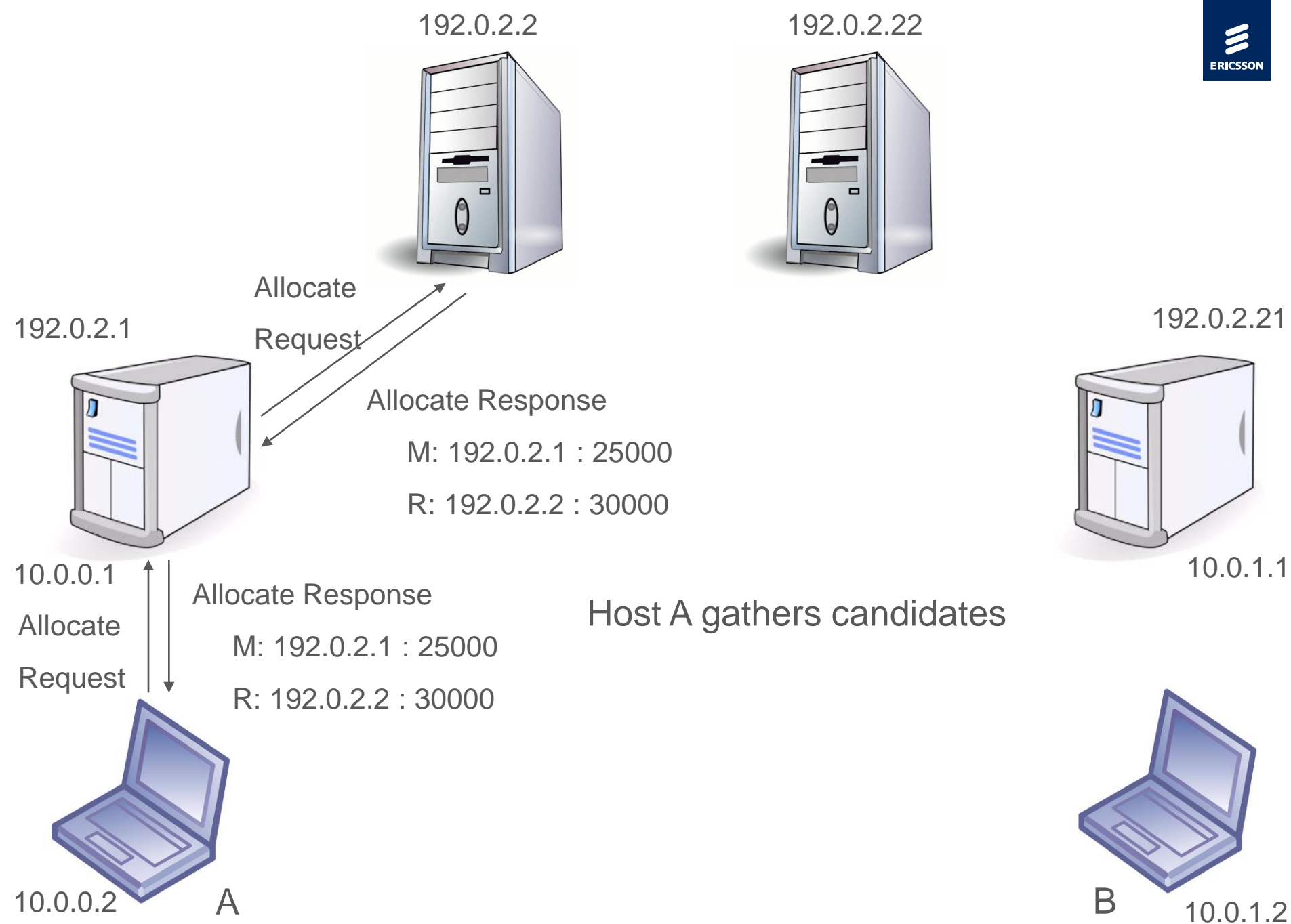


192.0.2.23



ICE Example (2)

- › Both endpoints are behind NATs
- › Endpoints use TURN servers



Host candidate:
10.0.0.2 : 20000
Server reflexive:
192.0.2.1 : 25000
Relayed:
192.0.2.2 : 30000

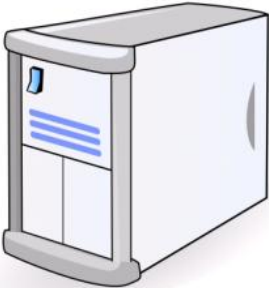
192.0.2.2



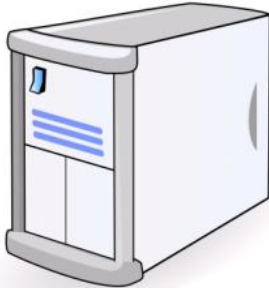
192.0.2.22



192.0.2.1



192.0.2.21



10.0.0.1

Allocate Response

M: 192.0.2.1 : 25000
R: 192.0.2.2 : 30000

... and forms a candidate list
that is sent to host B

10.0.1.1



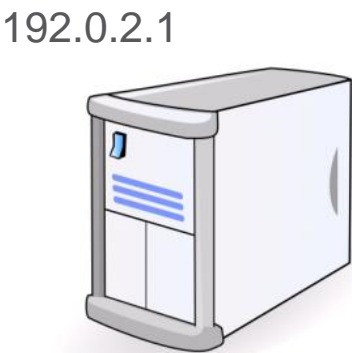
10.0.0.2

INVITE (offer)



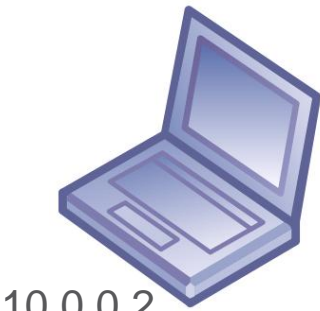
10.0.1.2

Host candidate:
 10.0.0.2 : 20000
 Server reflexive:
 192.0.2.1 : 25000
 Relayed:
 192.0.2.2 : 30000



10.0.0.1

Host B gathers candidates



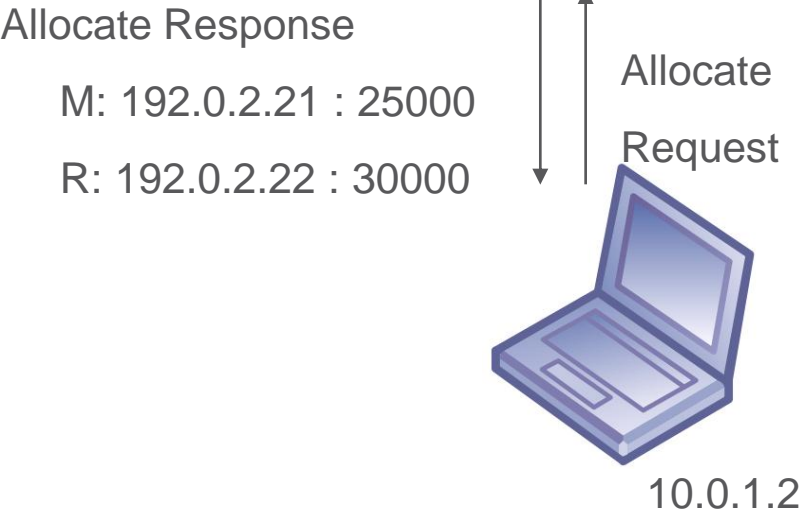
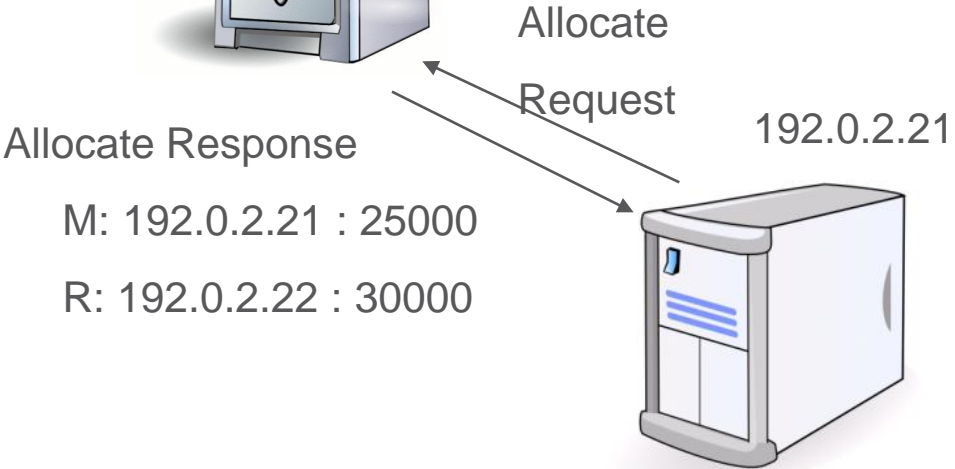
10.0.0.2



192.0.2.2



192.0.2.22



Host candidate:
 10.0.0.2 : 20000
 Server reflexive:
 192.0.2.1 : 25000
 Relayed:
 192.0.2.2 : 30000



Host candidate:
 10.0.1.2 : 20000
 Server reflexive:
 192.0.2.21 : 25000
 Relayed:
 192.0.2.22 : 30000



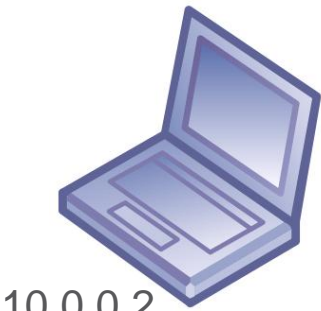
10.0.0.1

10.0.1.1

Allocate Response

M: 192.0.2.21 : 25000
 R: 192.0.2.22 : 30000

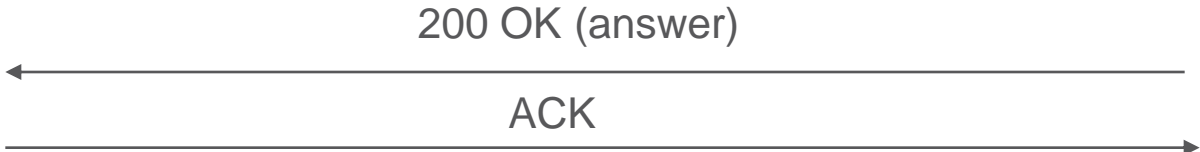
... and sends them to host A



10.0.0.2



10.0.1.2



Host candidate:
10.0.0.2 : 20000
Server reflexive:
192.0.2.1 : 25000
Relayed:
192.0.2.2 : 30000



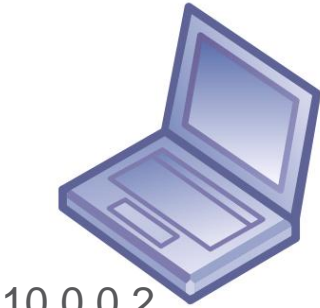
Host candidate:
10.0.1.2 : 20000
Server reflexive:
192.0.2.21 : 25000
Relayed:
192.0.2.22 : 30000



10.0.0.1

10.0.1.1

Connectivity checks sent to
host candidates fail due to
hosts being in different subnets



10.0.0.2

Binding Request



Packets Dropped

Binding Request



10.0.1.2

192.0.2.2



192.0.2.22



Host candidate:

10.0.1.2 : 20000

Server reflexive:

192.0.2.21 : 25000

Relayed:

192.0.2.22 : 30000

Host candidate:

10.0.0.2 : 20000

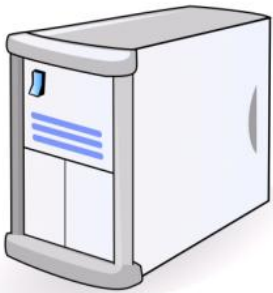
Server reflexive:

192.0.2.1 : 25000

Relayed:

192.0.2.2 : 30000

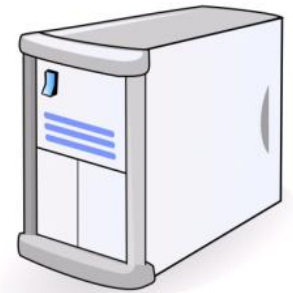
192.0.2.1



Binding Request

Packet Dropped

192.0.2.21



10.0.1.1

10.0.0.1

Binding

Request



10.0.0.2

B's NAT implements address dependent filtering



10.0.1.2

Host candidate:
 10.0.0.2 : 20000
 Server reflexive:
 192.0.2.1 : 25000
 Relayed:
 192.0.2.2 : 30000

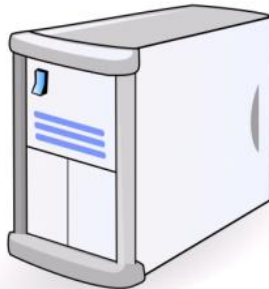


Host candidate:
 10.0.1.2 : 20000
 Server reflexive:
 192.0.2.21 : 25000
 Relayed:
 192.0.2.22 : 30000

192.0.2.1



192.0.2.21



Binding Request

Binding Response

10.0.0.1

Binding
Request

Binding Response



10.0.0.2

Also A's NAT implements address dependent filtering, but has now a binding for B's mapped address (due to the earlier connectivity check)

Binding
Request

Resp.



10.0.1.2

Host candidate:
 10.0.0.2 : 20000
 Server reflexive:
 192.0.2.1 : 25000
 Relayed:
 192.0.2.2 : 30000

192.0.2.2



192.0.2.22

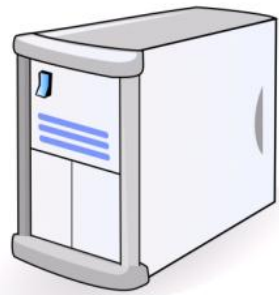


Host candidate:
 10.0.1.2 : 20000
 Server reflexive:
 192.0.2.21 : 25000
 Relayed:
 192.0.2.22 : 30000

192.0.2.1



192.0.2.21



Binding Response

Binding Request

10.0.0.1

Binding
 Resp.

Binding Request



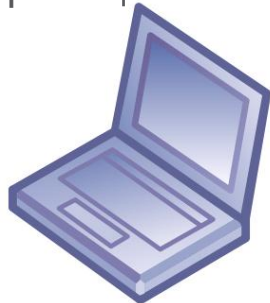
10.0.0.2

A performs a triggered check
 which now succeeds (there is
 a binding in B's NAT too)

10.0.1.1

Binding
 Resp.

Req.



10.0.1.2

Host candidate:
10.0.0.2 : 20000
Server reflexive:
192.0.2.1 : 25000
Relayed:
192.0.2.2 : 30000

192.0.2.2

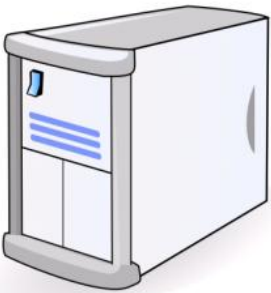


192.0.2.22



Host candidate:
10.0.1.2 : 20000
Server reflexive:
192.0.2.21 : 25000
Relayed:
192.0.2.22 : 30000

192.0.2.1

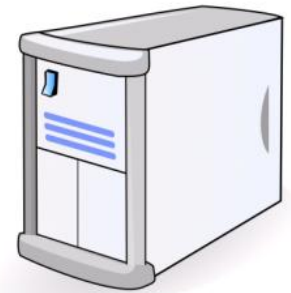


Binding
Request

Binding Response

Data
Indication

192.0.2.21



Send
Indication

10.0.0.1

Binding
Request

Binding Response



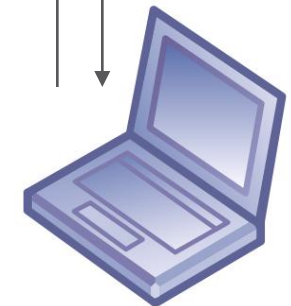
10.0.0.2

Further checks may be done
until stopping criteria is met

Send
Indication

Data
Indication

10.0.1.1



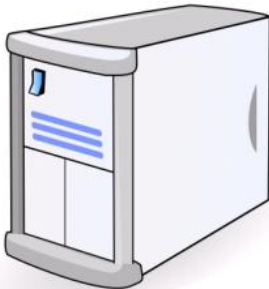
10.0.1.2

Host candidate:
 10.0.0.2 : 20000
 Server reflexive:
 192.0.2.1 : 25000
 Relayed:
 192.0.2.2 : 30000

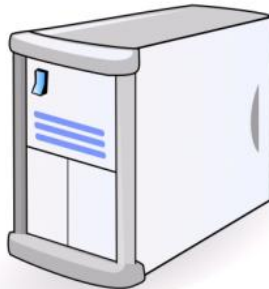


Host candidate:
 10.0.1.2 : 20000
 Server reflexive:
 192.0.2.21 : 25000
 Relayed:
 192.0.2.22 : 30000

192.0.2.1



192.0.2.21



Binding Response

Binding Request
 USE-CANDIDATE

10.0.0.1

Binding
 Resp.

Binding Request
 USE-CANDIDATE



10.0.0.2

Finally, controlling agent
 nominates the highest priority
 pair for use

10.0.1.1

Binding
 Resp.

Req.



10.0.1.2

ICE TCP

- › TCP Candidates with Interactive Connectivity Establishment (draft-ietf-mmusic-ice-tcp-09)
- › Establishing TCP connections with ICE
- › Somewhat low success ratio compared to UDP case
- › Uses active, passive and TCP simultaneous-open candidates
- › Recommends to use also any other means available

Other NAT Traversal Methods

- › Middle box communications
 - Signaling with NATs to create proper state in them
 - UPnP, SOCKS, MIDCOM, etc.
- › UDP/TCP hole punching
 - Number of variations for creating NAT bindings by sending packets to different addresses
 - One of the techniques used by ICE
- › Transparently for applications
 - Teredo (own variant of UDP hole punching and IPv6 over UDP)
 - Host Identity Protocol (uses ICE and UDP encapsulation)
- › ...

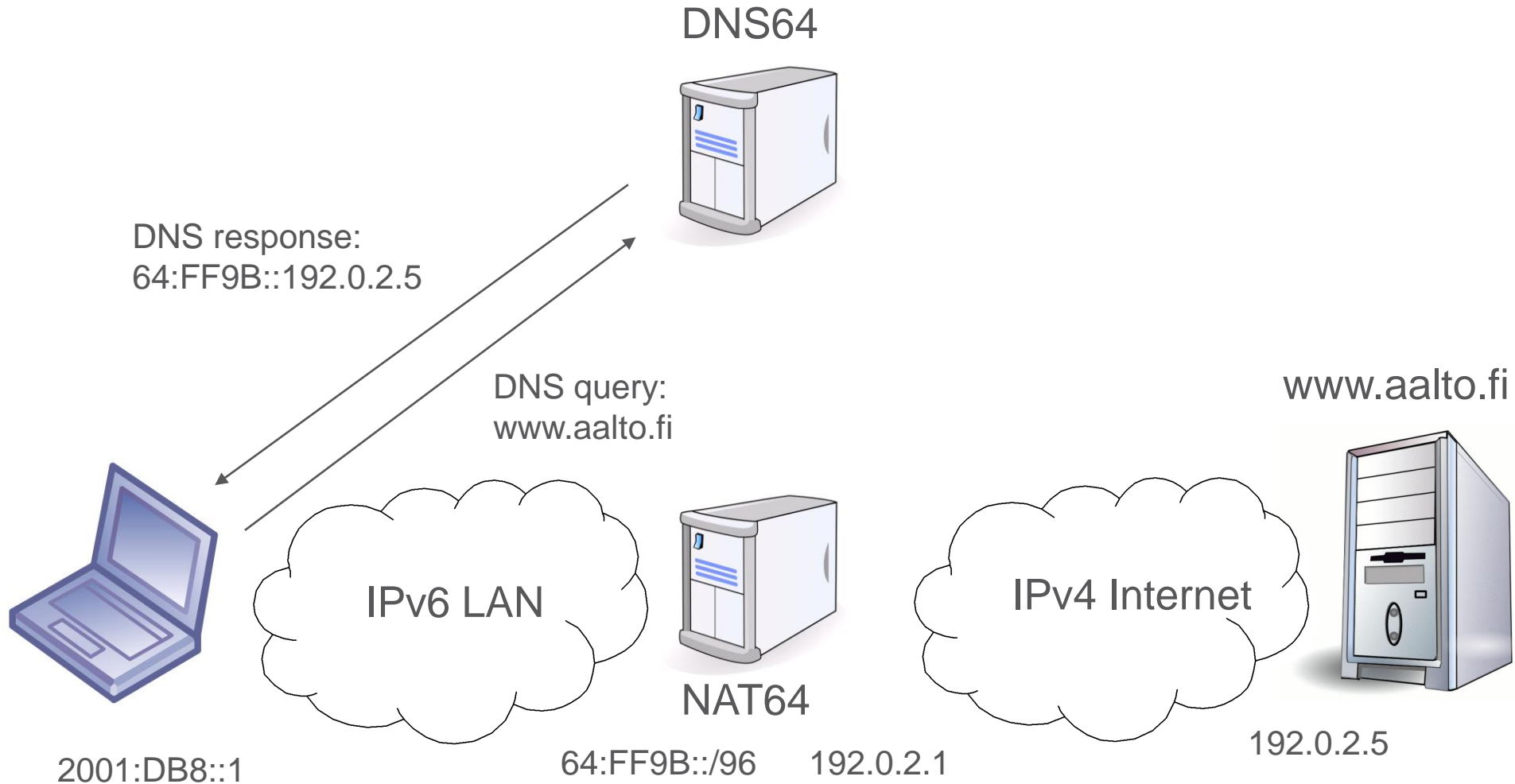
Outline

- › Introduction to NATs
- › NAT Behavior
 - UDP
 - TCP
- › NAT Traversal
 - STUN
 - TURN
 - ICE
 - Others
- › NAT64

NAT64 and DNS64

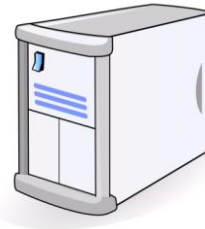
- › A client in IPv6-only network may need to communicate with a server in the IPv4-Internet
- › NAT64 translates packets between IPv6 and IPv4
- › DNS64 generates IPv6 addresses for servers that do not have one
 - Uses specific IPv6-prefix for routing traffic via the NAT64
 - Problems with hosts without a DNS entry

DNS64

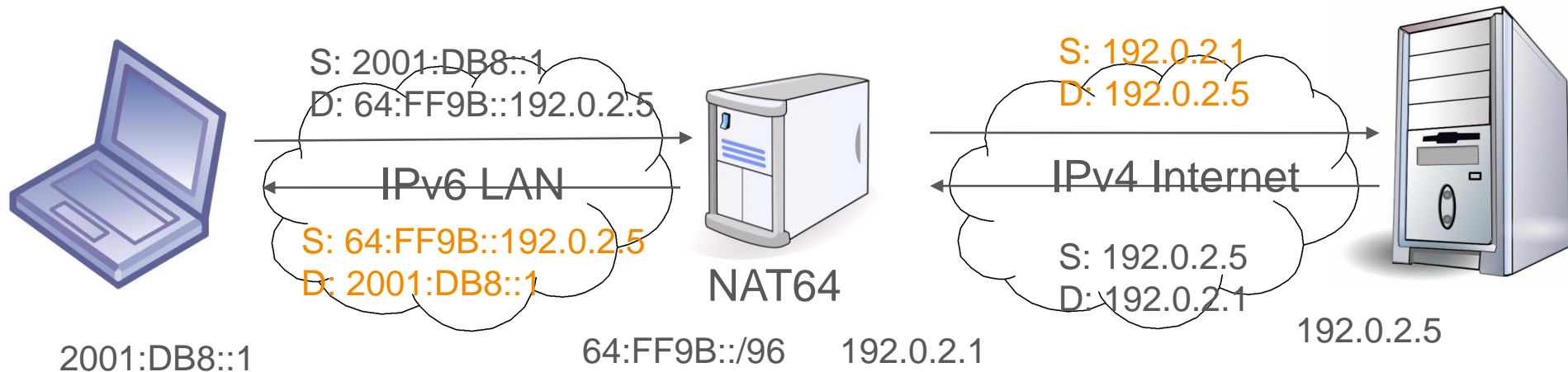


NAT64

DNS64



www.aalto.fi



Summary

- › NA(P)Ts originally invented to save IPv4 addresses
 - Can serve a whole subnet with a single IP address
 - Works (fairly well) for client-server, but breaks P2P connectivity
- › NATs have different (and often un-deterministic) behavior
 - Endpoint-(in)dependent mapping and/or filtering
 - IP address and port assignment, timeouts, etc.
- › NAT traversal developed to fix connectivity
 - STUN and TURN for server-reflexive and relayed addresses
 - ICE uses STUN and TURN for gathering candidates and running connectivity checks between them; tries various possible combinations and selects the best
- › NAT64 provides IPv4 connectivity when network only provides IPv6